



**MSTP**

# INTRODUCTION TO NETWORKING ESSENTIALS AND TCP/IP



# TCP/IP Course Outline

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**MSTP**

- OSI Reference Model
- LAN Segmentation
- Internet Protocol Stack
- IP Addressing and Subnetting



# Open Systems

## Interconnect (OSI)

**MSTP**

- OSI is a Layered Network Model for networking protocols. Establishes standards for internetworking.
- Allows for shortcut explanations to facilitate protocol comparisons.
- OSI is not a popular method for interconnecting computers. WHY? Seven layer stack is extremely complicated and next to impossible to keep track of for standardization purposes.



# Why a layered Network Model?

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- Clarify what general functions are to be done rather than how to do it.
- Reduce the complexity of networking into more manageable sub-layers.
- Enable interoperability using standard interfaces (Application Programmable Interface).
- Allow changes in one layer to occur without changing other layers.
- Allows specialization within the network industry.



# Physical and Logical Data Movement

**MSTP**

- Physical movement of data
  - Application layer protocol
    - Someone creates information on an application.
  - Communication protocol
    - The information is then packaged for transmission.
  - Transmission protocol
    - The package is now prepared for actual physical transmission.
- Logical movement of data
  - Physical topology
    - The data moves across some type of physical channel.





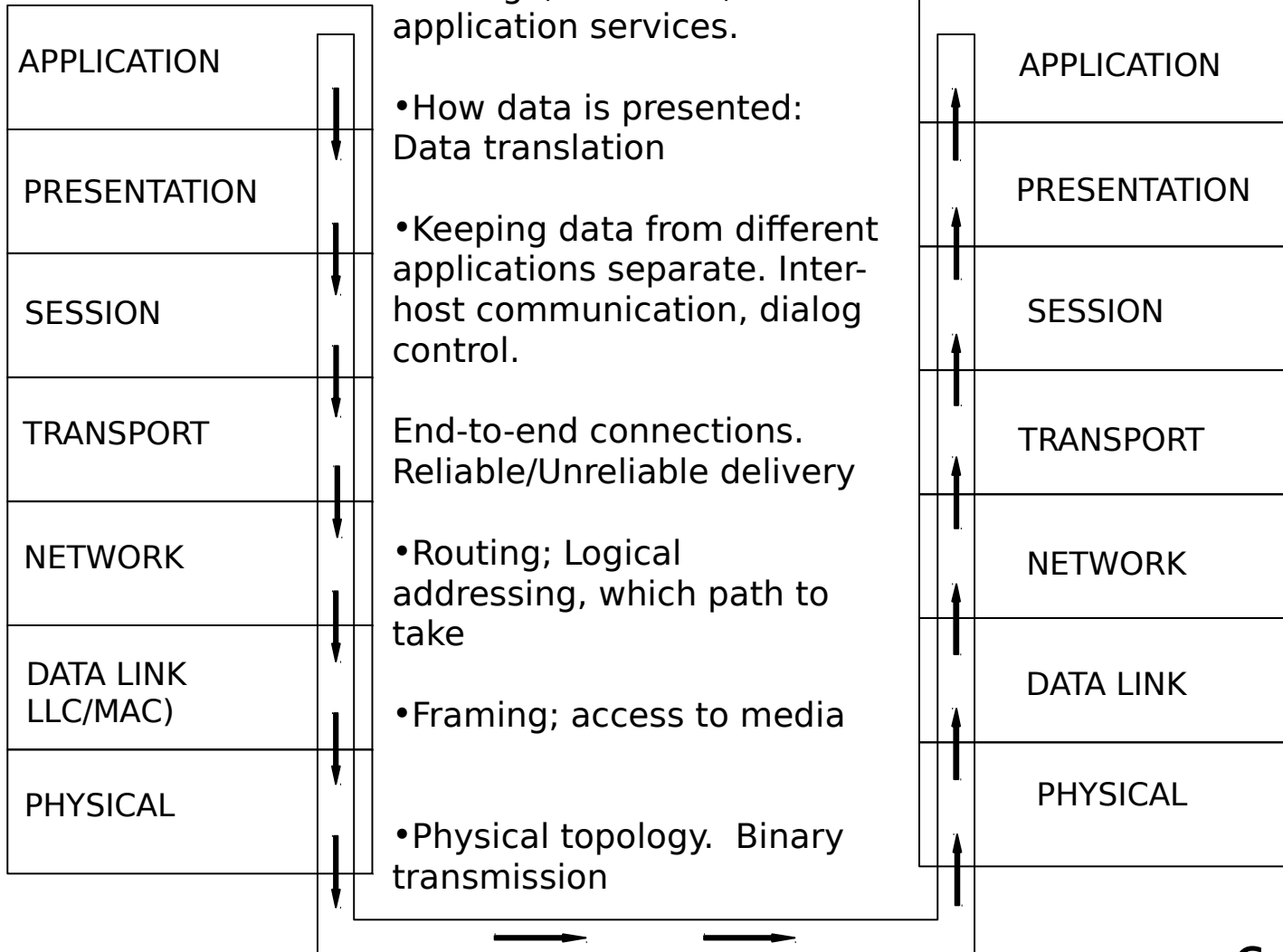
# OSI Protocol Stack

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Application layer protocol

Communication layer protocol

Physical topology





# Application Layer

**MSTP**

APPLICATION
PRESENTATION
SESSION
TRANSPORT
NETWORK
DATA LINK LLC/MAC)
PHYSICAL

- The application layer identifies and establishes the availability of intended communication partners.
- Synchronizes cooperating applications.
- Establishes agreement on procedures for error recovery and control of data integrity.



# Application Layer cont...

**MSTP**

## Network Applications

- Electronic Mail
- File Transfer
- Remote Access
- Client/Server Process
- Network Management
- Others

## Internetwork Applications

- Electronic Data Interchange
- World Wide Web
- E-Mail Gateways
- Special-Interest Bulletin Boards
- Financial Transaction Services

Internet Navigation Utilities

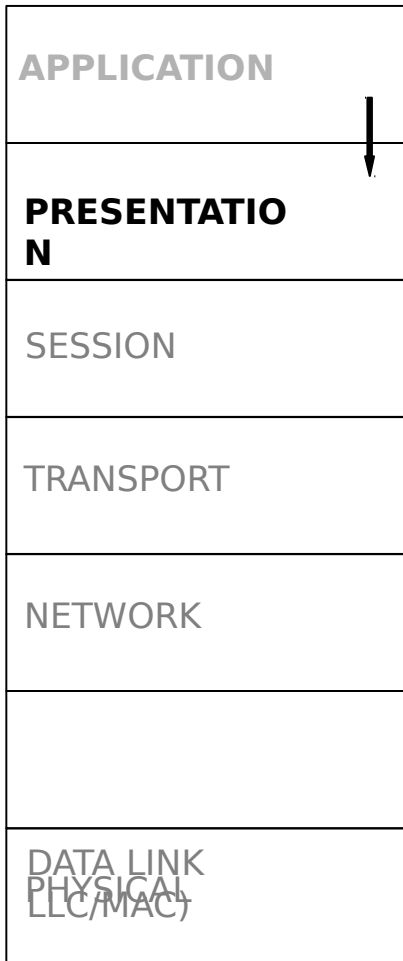
- Conferencing (Voice, Video, & Data)





# Presentation Layer

**MSTP**

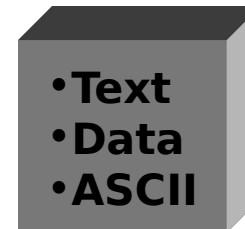
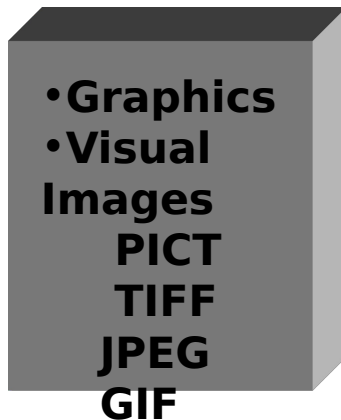


- This layer ensures that information sent by the application layer of one system will be readable by the application layer of another.
  - Data translation
  - Encryption
  - Compression
- Negotiates data transfer syntax for the application layer.



# Presentation Layer

**MSTP**



Capt  
Gaughen

**Presentation Layer provides code  
conversion**

# Presentation Layer (Cont....)



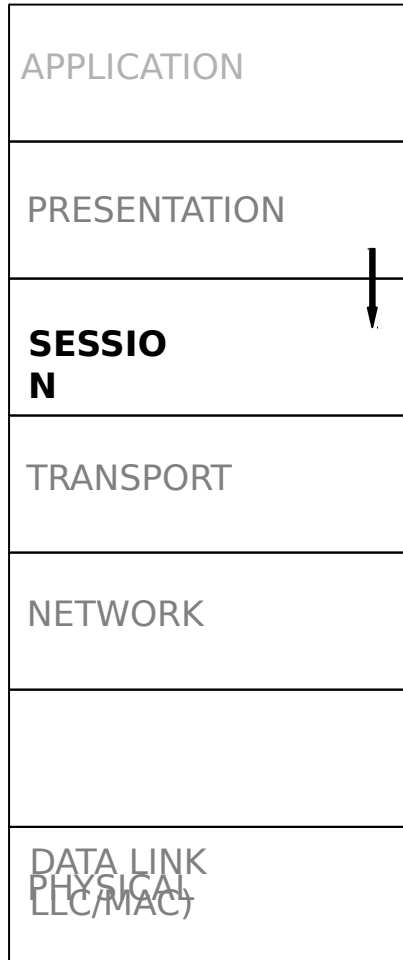
**MSTP**





# Session Layer

**MSTP**



- This layer establishes, manages, and terminates sessions between applications by offering three modes:
  - Simplex (monologue)
  - Half-duplex (forbidden interruption)
  - Full-duplex (flow control issue)
- Accomplished in three phases:
  - Connection establishment
  - Data transfer
  - Connection release
- Manages data exchange between presentation layer and entities.



# Half-duplex and Full-duplex

**MSTP**

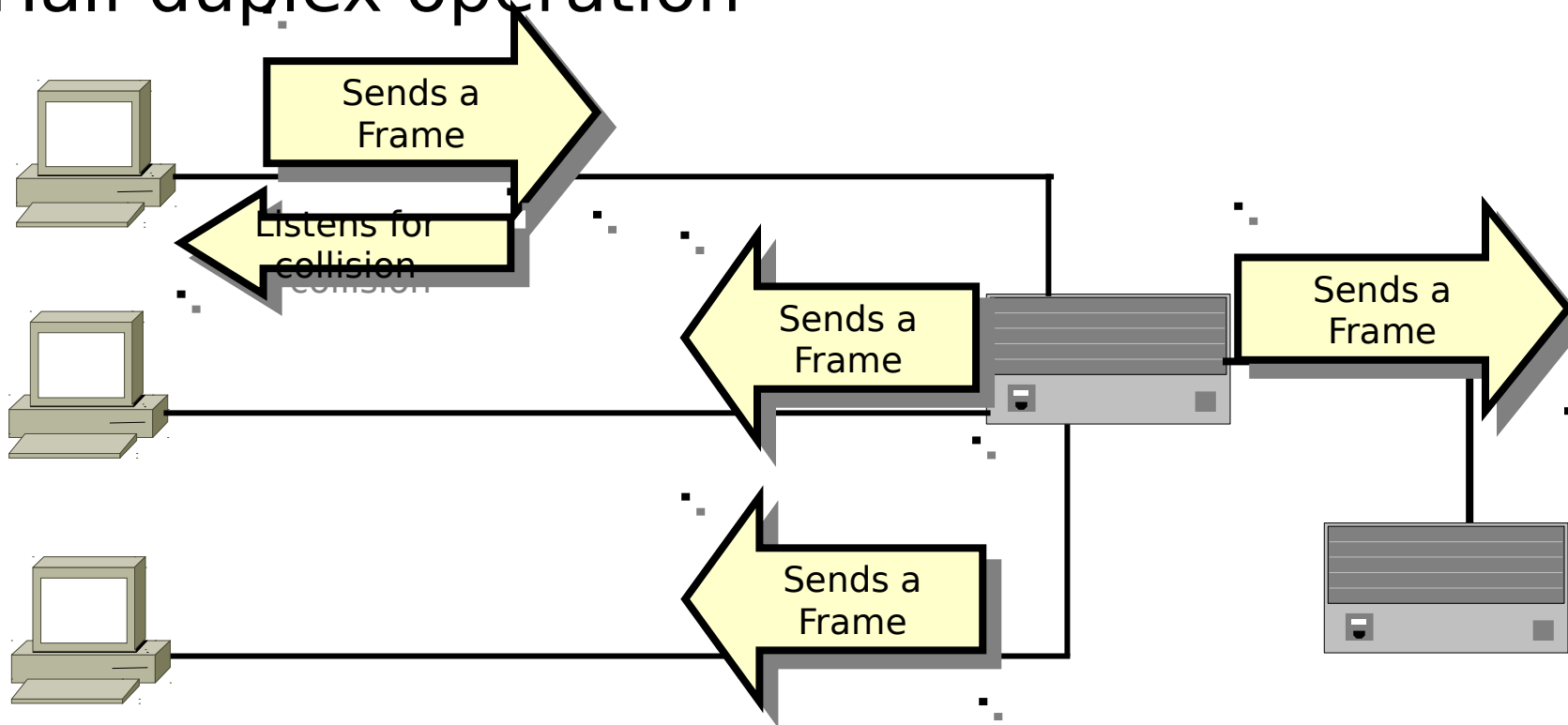
- Half-duplex – Nodes take turns transmitting and receiving. Ex. One way bridge
- Full-duplex – Nodes can transmit and receive simultaneously, but it requires a switch port, not a hub, to be able to do so.



# Half Duplex

**MSTP**

- Half-duplex operation

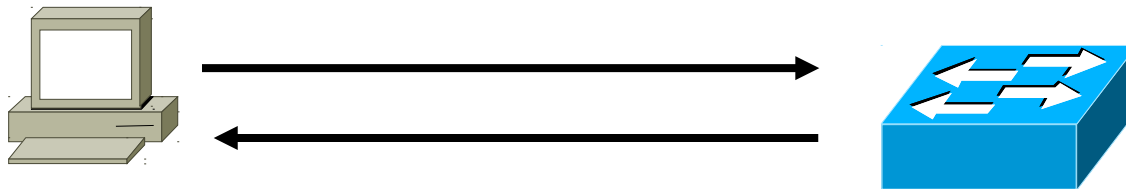




# Full Duplex

**MSTP**

- Full duplex is allowed when the possibility of collision is removed.



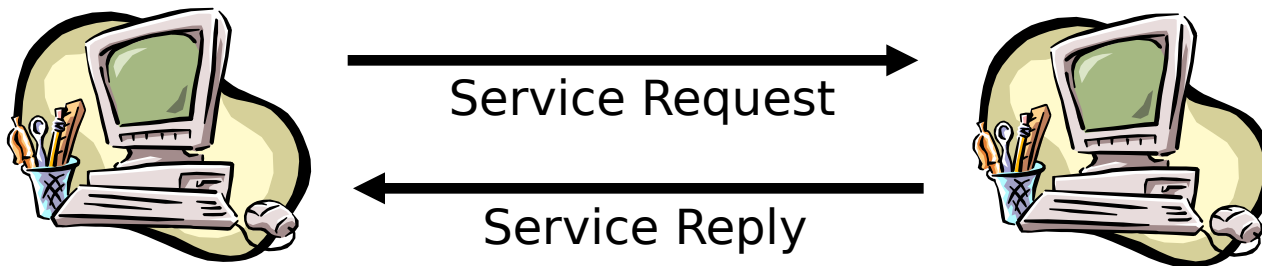
Because no collisions are possible, both ends can send and receive simultaneously. This reduces Ethernet congestion which results in low latency, information is moving in



# Session-Layer Protocols and Interfaces

**MSTP**

- Network File System (NFS)
- Structured Query Language (SQL)
- Remote-Procedure Call (RPC)
- X Window System
- NetBios Names
- Internet Browsers
- DNA Session Control Protocol (SCP)

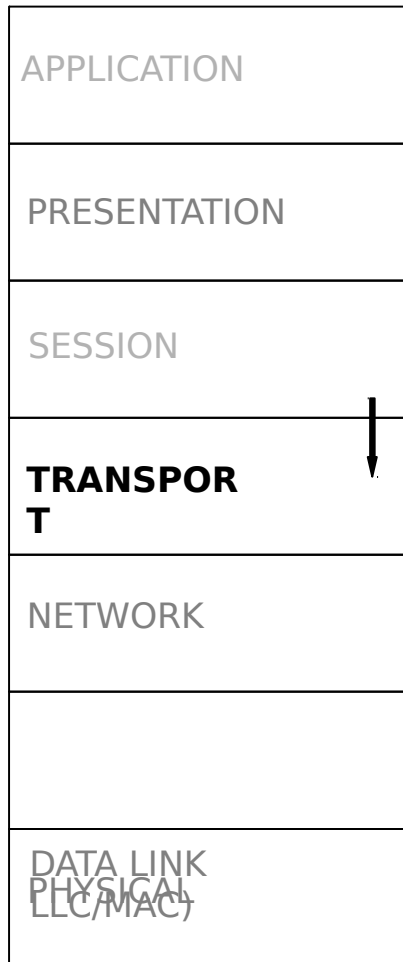






# Transport Layer

**MSTP**



- Reliable network communication between end nodes
- Provides mechanisms for the establishment, maintenance, and termination of virtual circuits.
- Transport fault detection and recovery
- Information flow control (buffering, windowing, congestion avoidance)



# Buffering

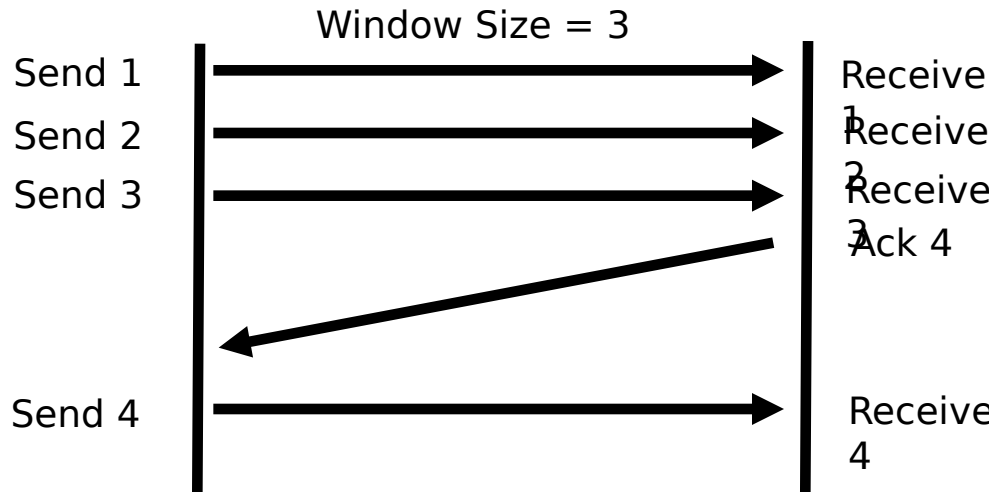
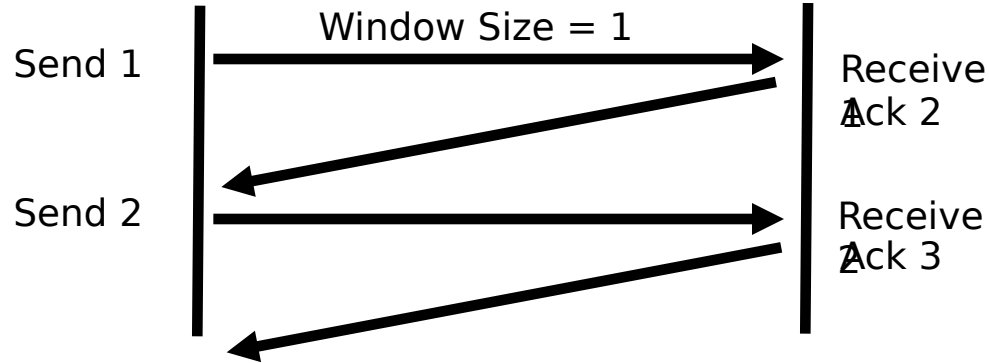
**MSTP**

- Computers reserve enough buffer space that bursts of incoming data can be held until processed.
- Cannot slow down the transmission rate of the sender that is sending the data.
- Common method of dealing with the rate of arrival of data



# Windowing

**MSTP**



Windows: 8,190  
bytes  
Cisco: 4,092 bytes



# Congestion Avoidance

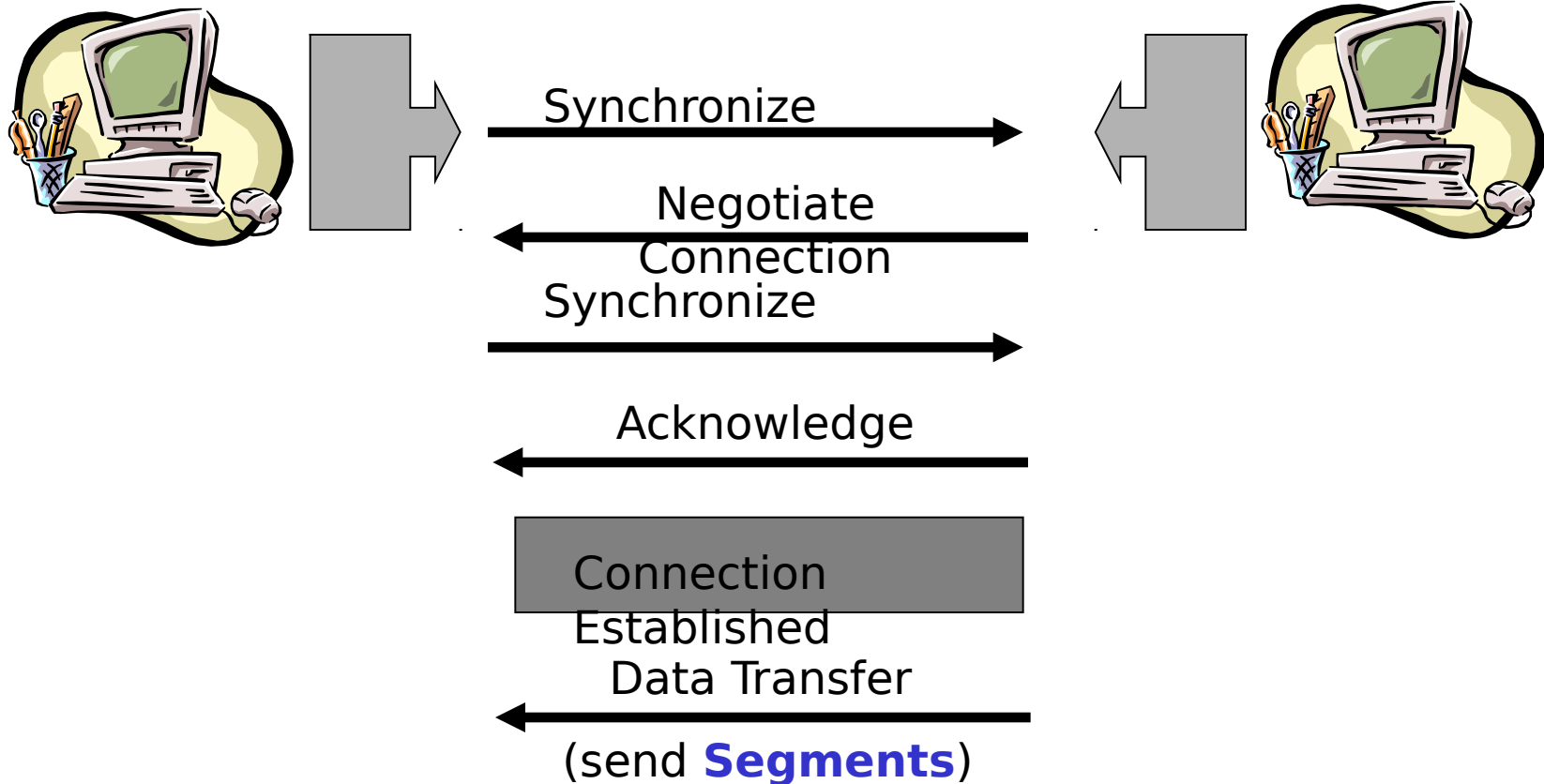
**MSTP**

- When a computer sees that its buffers are getting full, it sends a message to the sender to slow down sending messages.
- This process is used by Synchronous Data Link Control (SDLC) and Link Access Procedure, Balanced (LAPB) serial data link protocols.
- Ex. TCP/IP Internet Control Message Protocol (ICMP) message “Source Quench”
  - Sent by the receiver or some intermediate router to slow the sender. The sender will slow down gradually until “Source Quench” messages are no longer received.



# Setting Up a Reliable Session

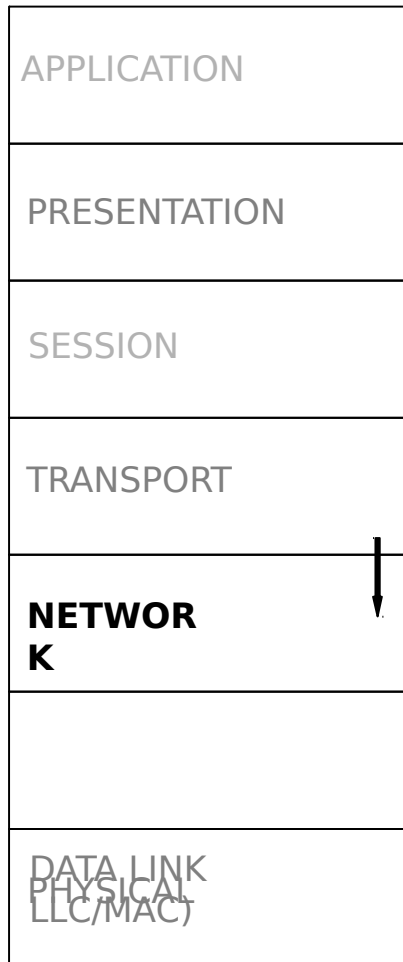
**MSTP**





# Network Layer

**MSTP**

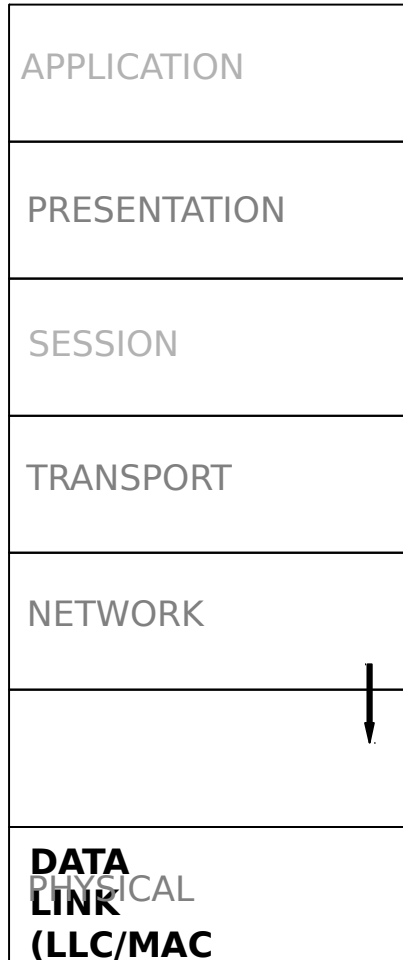


- The network layer is the layer at which routing occurs.
- This layer provides connectivity and path selection between two end systems.



# Data Link

**MSTP**

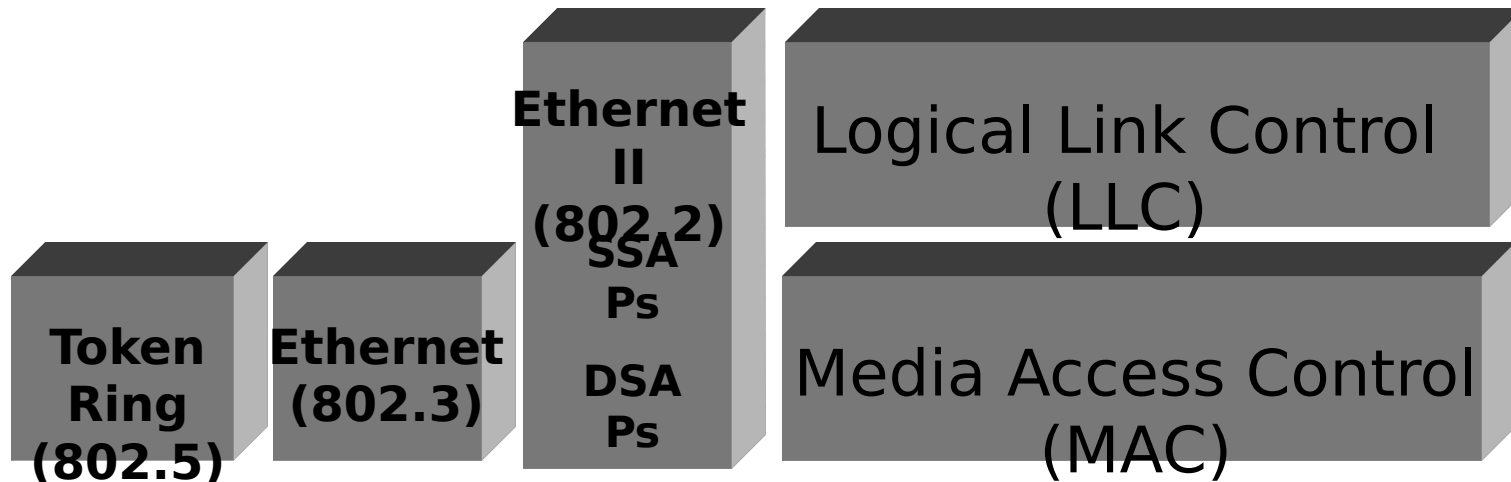


- This layer provides reliable transit of data across a physical link.
- Physical addressing (MAC)
- Network topology
- Line discipline (CSMA/CD & CSMA/CA)
- Error notification
- Divided into two sub-layers (MAC and LLC)



# Data Link Sub-Layers

**MSTP**



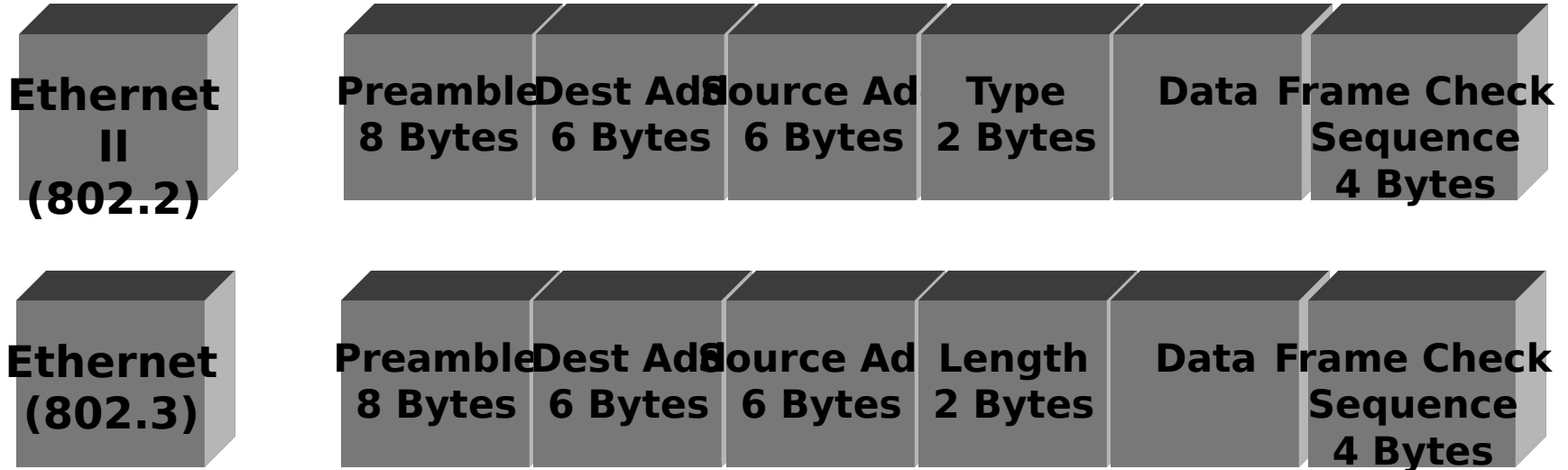
- Logical Link Control acts as the managing buffer
  - Source Service Access Points (SSAPs)
  - Destination Service Access Points (DSAPs)





# Data Link Sub-Layers

**MSTP**





# Physical Layer

**MSTP**

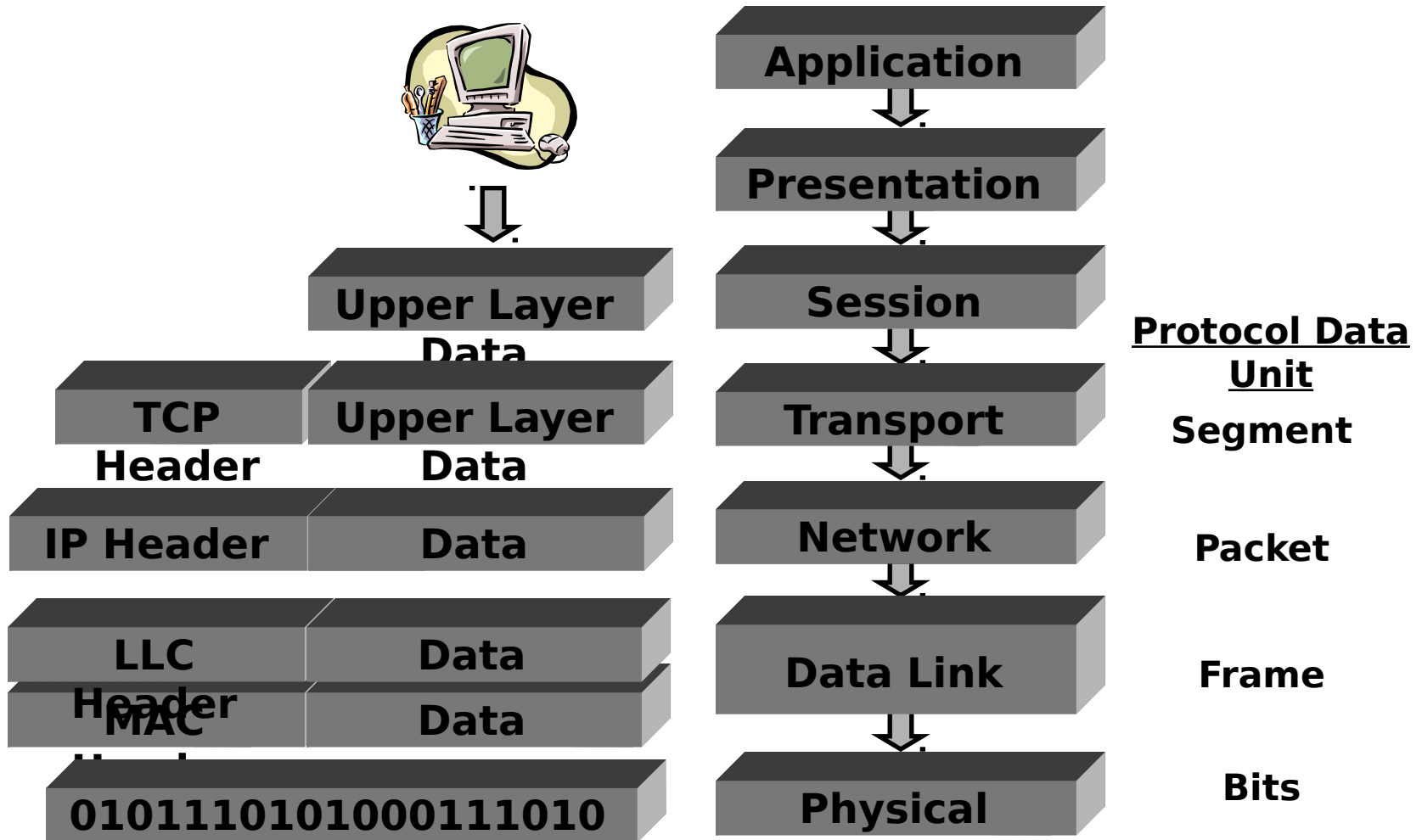
APPLICATION
PRESENTATION
SESSION
TRANSPORT
NETWORK
DATA LINK PHYSICAL (LLC/MAC)

- The physical layer defines the electrical, mechanical, procedural and functional specifications for activating, maintaining, and deactivating the physical link between end systems.



# Putting it all together ...

**MSTP**

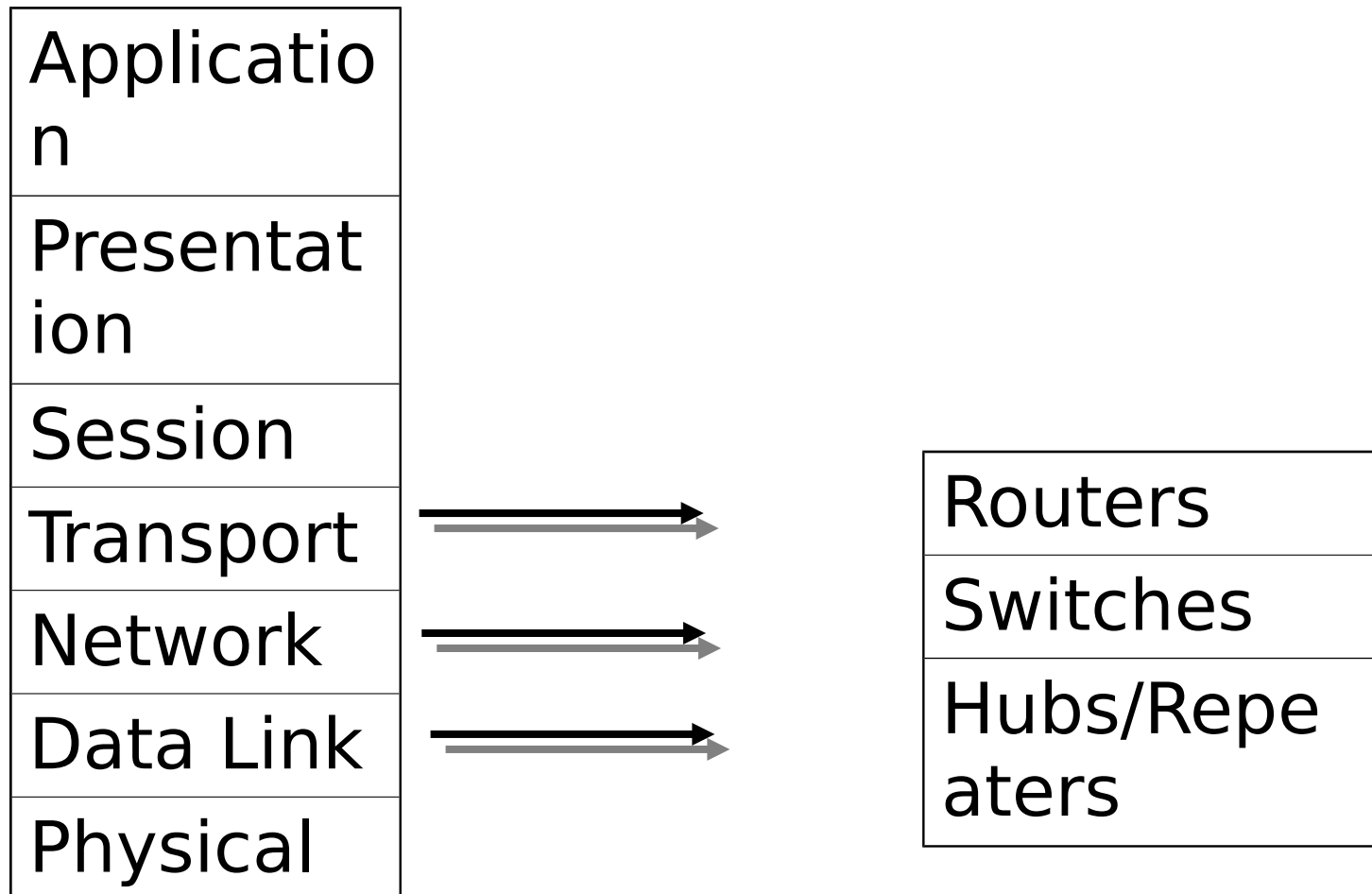




# Devices at the different layers

**MSTP**

- OSI Model





**MSTP**

# **LAN Segmentation & Networking Essentials**



# Internetworking Fundamentals

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- Internetworks are the communication structures that work to tie Local Area Networks (LAN) and Wide Area Networks (WAN) together.
- Primary goal is to move information anywhere quickly upon demand and with complete integrity. Must be able to connect many different networks together to serve the organizations needs regardless of the type of physical media involved.



# Internetworking Devices

**MSTP**

- LANs were designed to operate in limited geographical areas, such as one floor of a building, or a single building.
- LANs connect PCs together so that they can access network resources.
- A LAN connects physically adjacent devices on the network media or cable.
  - LAN Devices include: Repeaters, Bridges, Hubs, Switches, Routers, and Gateways.



# Internetworking Devices

## Cont.

**MSTP**

- WANs extend beyond the LAN to connection networks located in different building, cities, states, and countries together.
- WANs are connected over ***serial*** lines.
  - WAN devices include: Routers, ATM Switches, X.25 and frame relay switches, modems, Channel Service Unit/Data Service Units (CSU/DSU), communication servers, and multiplexors.





# Network Congestion

**MSTP**

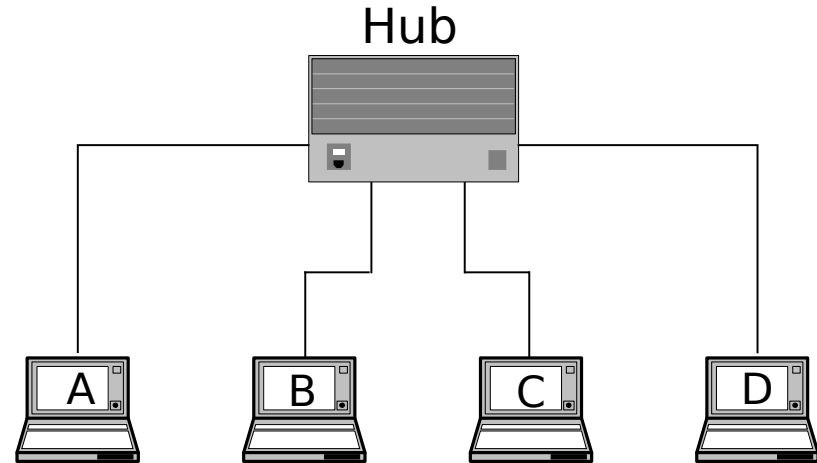
- Where does it come from?
- How do I get rid of it?
  - Physical segmentation
  - Bridges
  - Routers
  - Switches



# Hubs

**MSTP**

APPLICATION
PRESENTATION
SESSION
TRANSPORT
NETWORK
<b>Physical Layer</b>



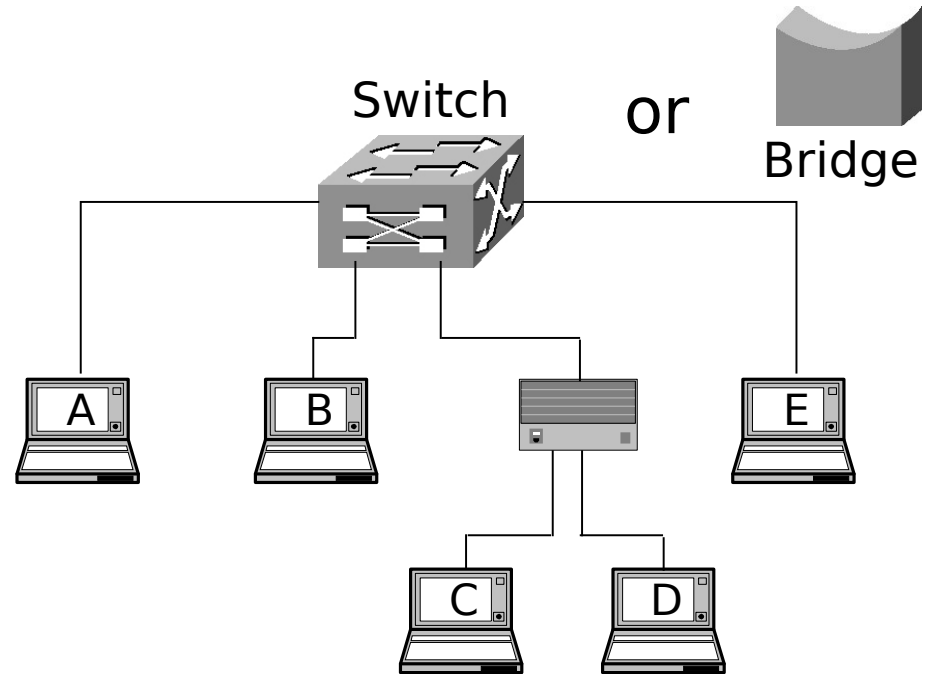
- All devices are in the same collision domain
- All devices are in the same broadcast domain
- All devices share the same bandwidth



# Switches/Bridges

**MSTP**

APPLICATION
PRESENTATION
SESSION
TRANSPORT
NETWORK
DATA LINK Physical LLC/MAC



- Each Segment has its own collision domain
- All segments are in the same broadcast domain
- Listening Learning Filtering and

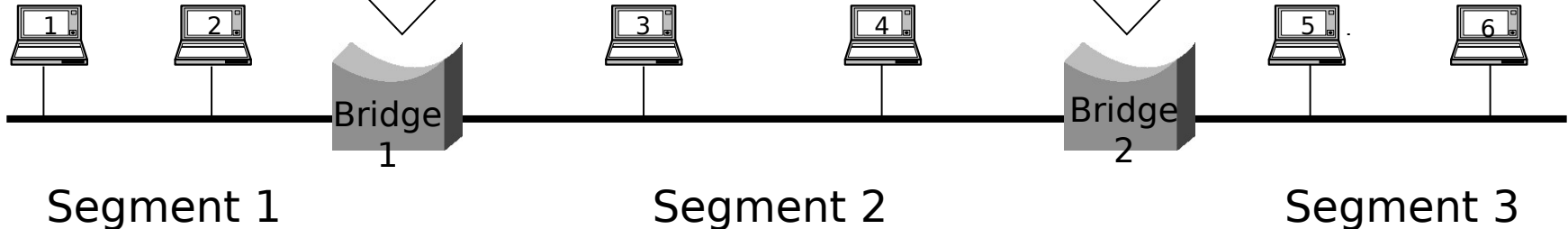


# Segmentation with Bridges

**MSTP**

Forwarding Table	
<u>Host</u>	<u>Segme nts</u>
1	1
2	1
3	2
4	2
5	2
6	2

Forwarding Table	
<u>Host</u>	<u>Segme nts</u>
1	1
2	2
3	2
4	2
5	3
6	3



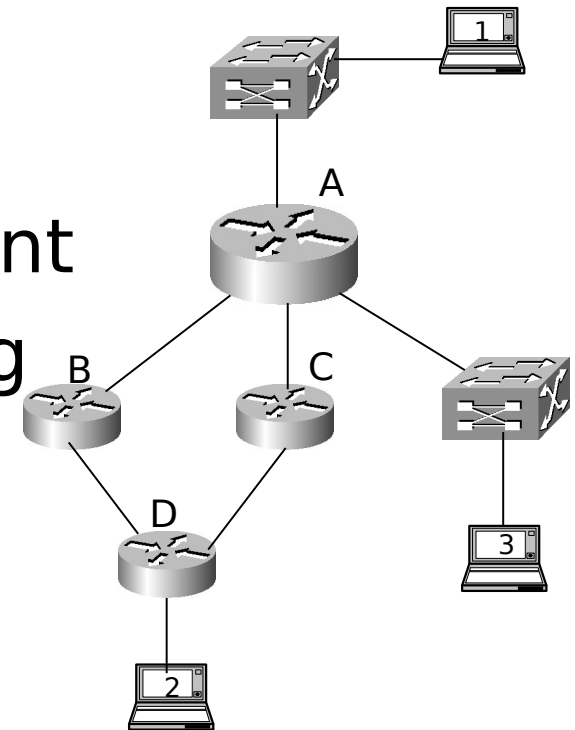


# Routers

**MSTP**

APPLICATION
PRESENTATION
SESSION
TRANSPORT
<b>NETWORK</b>
DATA LINK (Physical Layer) (MAC)

- Broadcast control
- Multicast control
- Optimal path determination
- Traffic Management
- Logical addressing
- Connects to WAN services





# Ethernet: Collisions

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- Certain level of collisions are expected on CSMA/CD LANs
- Excessive collisions can result from faulty components or overloaded segments
  - **Bad or excessively long cables**
  - **Bad NICs or transceivers**
- Establishing a baseline is helpful to determine normal levels
- Local collisions
  - **Occur on local LAN segment**
  - **Detected by circuitry in LAN interfaces**
- Remote collisions
  - **Occur on other side of repeater nodes**



# CSMA/CD

**MSTP**

1. Sender is ready to send the frame. It listens to detect whether any frame is currently being received.
2. If Ethernet is silent, the devices begins to send the frame.
3. The sending device begins to listen to ensure that the frame it is sending does not collide with a frame that another station is sending.
4. If no collision occurs, the bits of the sent frame are received back successfully.
5. If a collision occurred, the device sends a jam signal and then waits a random amount of time before repeating the process.



# Ethernet: Jabber

**MSTP**

- Jabber
  - Frames are longer than 1518 bytes
  - Fails CRC check
- Causes
  - Caused by faulty transceivers
  - Transceivers can transmit for only 150-millisecond intervals
    - Sufficient time to transmit 1518 bytes
  - If transceiver does not stop after 1518 bytes, jabber results





# Ethernet: Performance Issues

**MSTP**

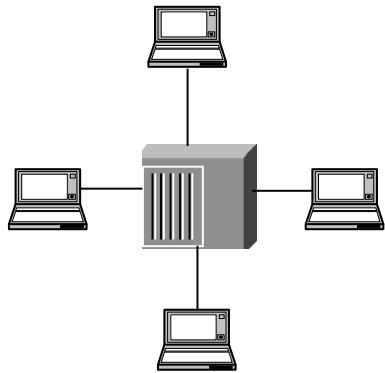
- Measuring network utilization
  - Protocol analyzers
  - User complaints
  - Rule of thumb: shared Ethernet segments < 40% utilization
- Improving network utilization
  - Segmenting with routers
  - Segmenting with switches



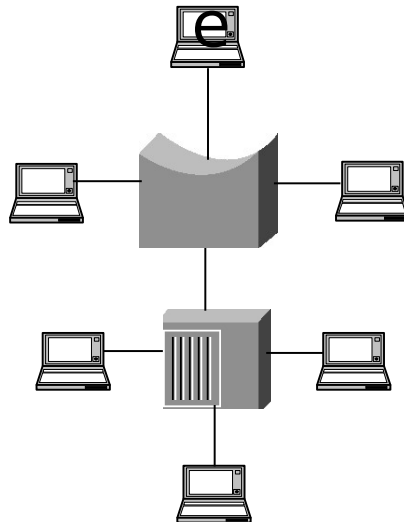
# Creating Domains

**MSTP**

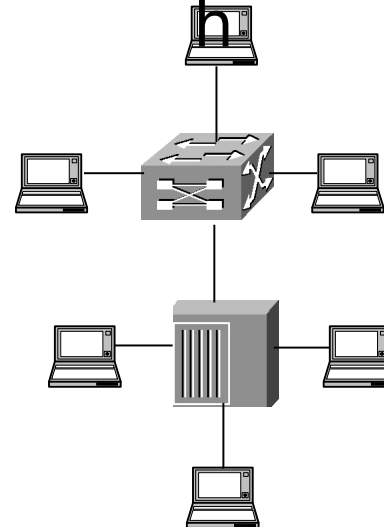
Hub



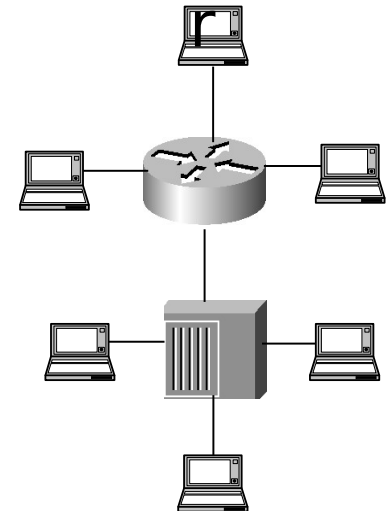
Bridge



Switch



Route



Collision

Domains:

Broadcasts

Domains:

1 4

4

1

4

4

1

42



**MSTP**

# **TCP/IP Protocol Stack**



# TCP/IP: What and Why?

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- **What is TCP/IP?**

- TCP/IP is the basic network protocol that is used as the base language for Joint Service and commercial worldwide communications . Think of it as the "body language" of network protocols, meaning that it is known worldwide as a basic means of communicating.

- **Why do we care about it?**

- Marine Corps tactical systems talk in native IP.



# What is TCP/IP...continued

**MSTP**

- TCP/IP is not a single protocol, but rather, a suite of protocols that work together to interconnect networks and provide a wide variety of services.
- Basic functions of TCP/IP are **remote login, file transfer, and email**. TCP/IP enables network devices to determine the physical address of LAN nodes, to map English language names to numeric machines names, and to manage the network.
- Protocols were created to work with virtually any host hardware, operating system and 45 connecting media. Networks built upon TCP/IP



# TCP/IP Timeline

**MSTP**

- ARPANET begins with 3 computers in California and one in Utah
- MID 1970's - DARPA researching packet switching technology. DARPA funded research of internetworking technology
- 1977 - 1979 - TCP/IP Standards defined (1969 - First work on protocol)
- 1980 - DARPA began converting research nets to TCP/IP

University of California (Berkeley) distributes BSD UNIX. DARPA funds UCB to integrate TCP/IP into UNIX TCP/IP on ARPANET and majority of college networks

- 1983 - OSD mandated TCP/IP for "all" long haul networks.

DOD adopts TCP/IP as standard. ARPANET divides



# TCP and IP

**MSTP**

- TCP
  - Connection Oriented
  - Requires Synchronization
    - Synchronization lost - reset session
  - Requires ACK
  - Sequencing numbers
    - No ACK - retransmit
    - Bad segment - retransmit all
  - Data Flow Control
    - Windowing
- IP:
  - Unreliable, “best effort”
    - No ACK required
    - No guarantee of delivery
      - Packets could get lost
      - Delivered out of sequence
      - Duplicated
      - Delayed



# OSI to TCP/IP Comparison

**MSTP**

**OSI**

**TCP/IP**

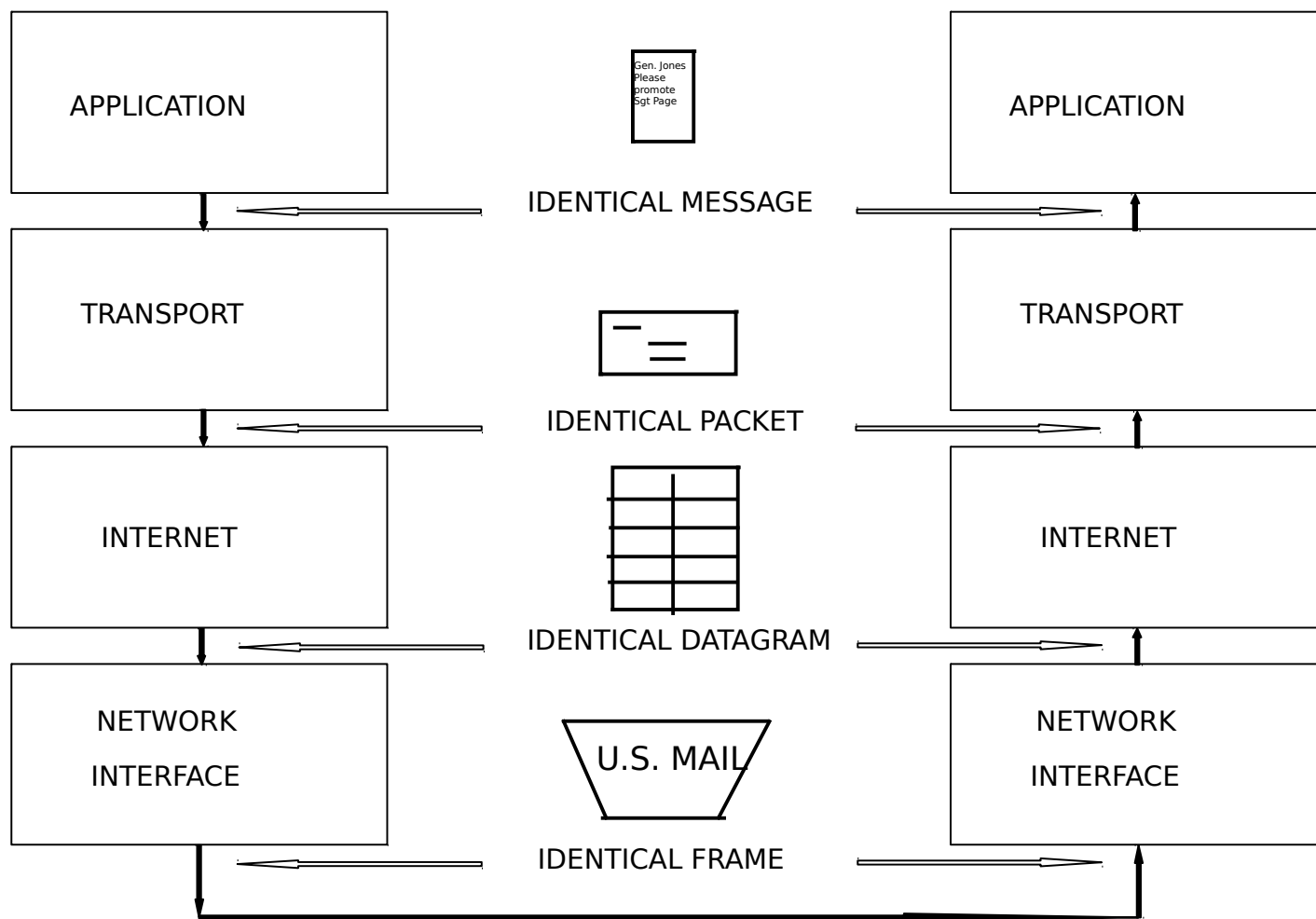
APPLICATION		PROCESS / APPLICATION
PRESENTATION		
SESSION		TRANSPORT (HOST-TO- HOST)
TRANSPORT		
NETWORK		INTERNET
DATA LINK		NETWORK ACCESS
PHYSICAL		





# TCP/IP Protocol Stack

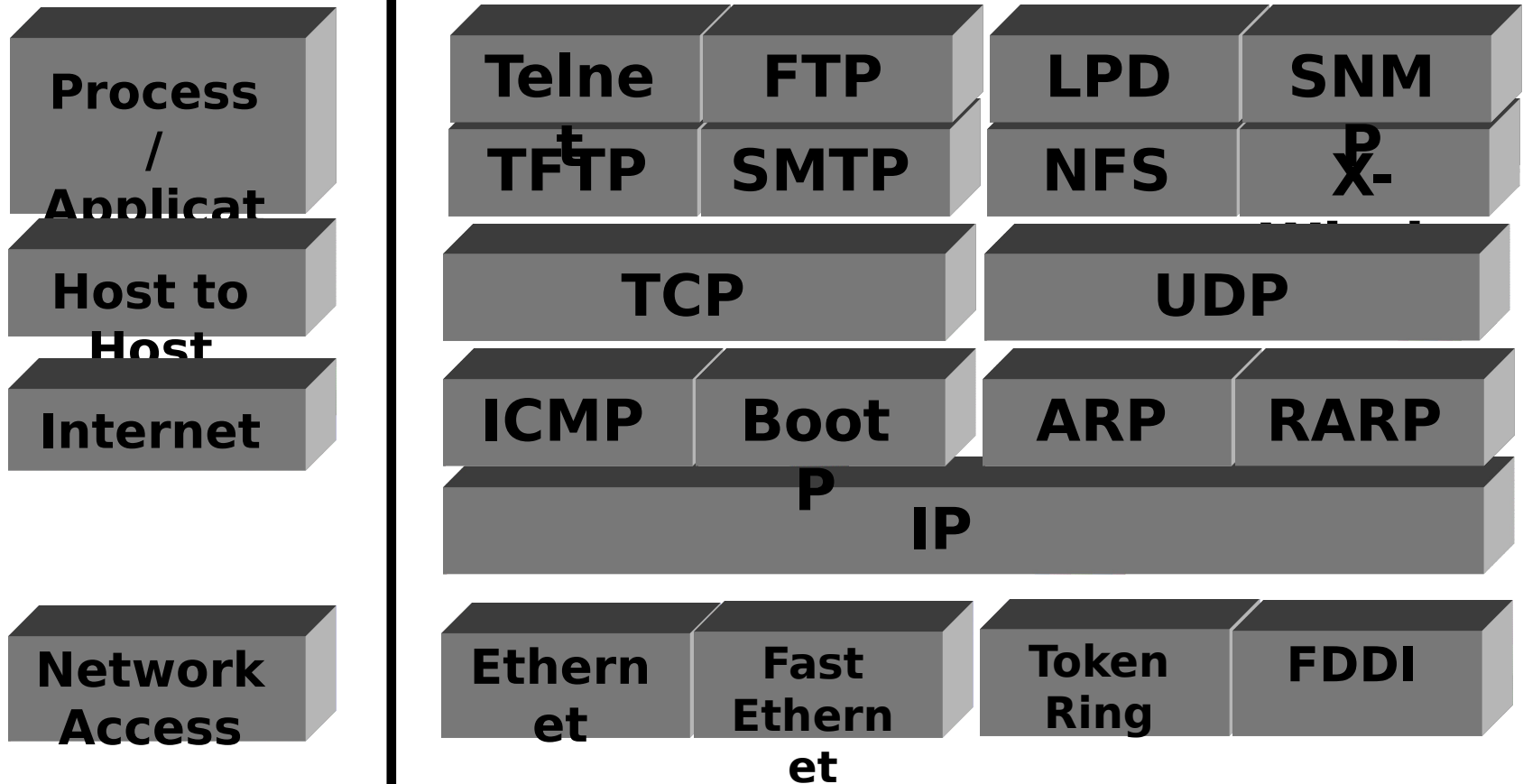
**MSTP**





# TCP/IP Protocol Stack

MSTP





# Transmission Control Protocol (TCP)

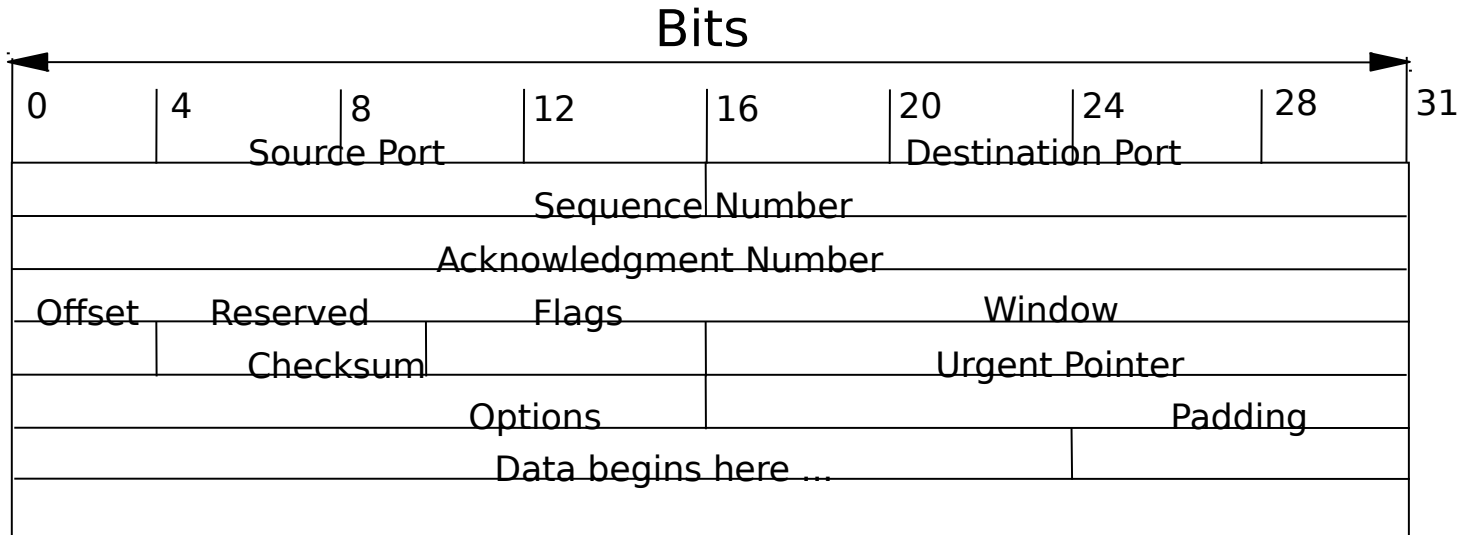
**MSTP**

- Connection oriented
  - Connection must be established prior to data transfer
  - Adds overhead
- In sequence delivery
  - Uses segment numbers to guarantee packet arrival in sequence deliver
  - Adds error checking & sequence numbering
- Provides graceful release
  - Ensures all data sent is received
- Reliable
  - Acknowledgment of received packets



# TCP Header

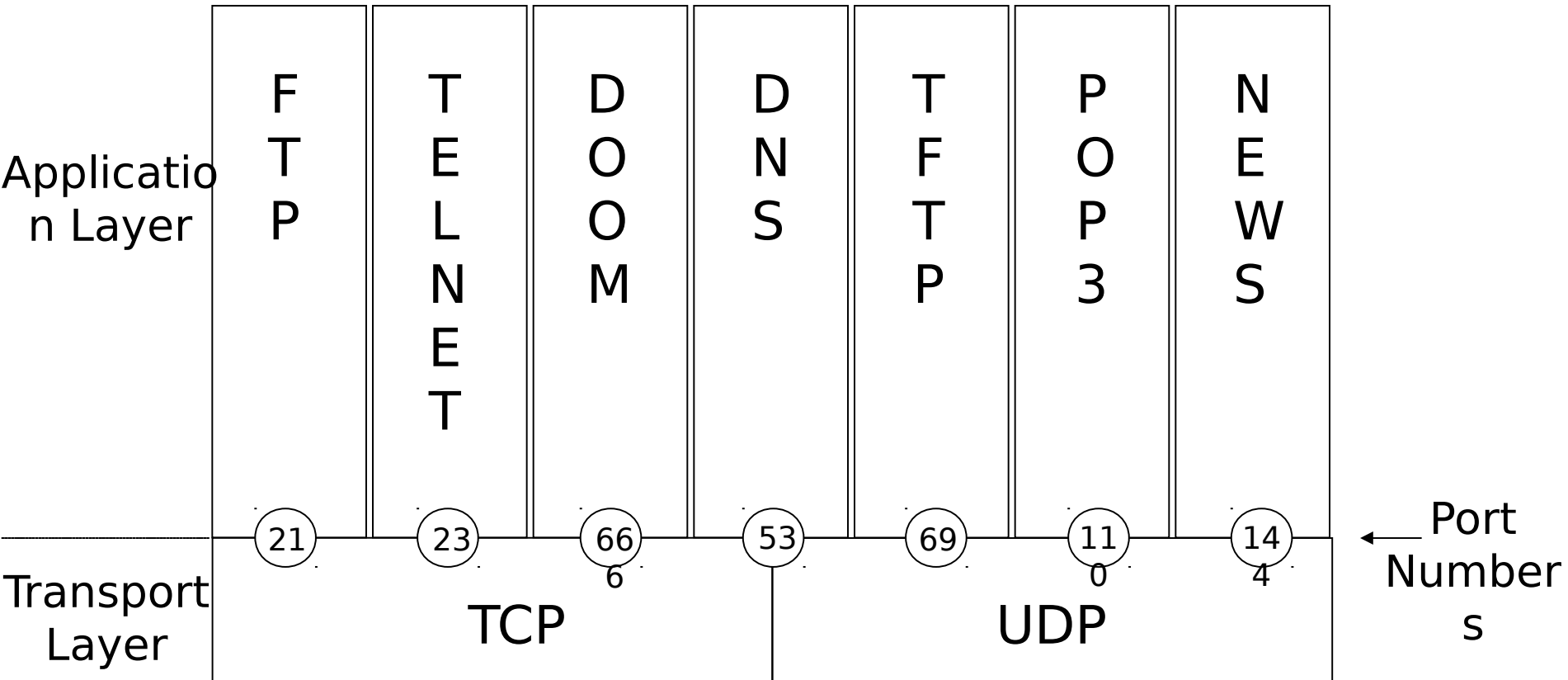
**MSTP**





# TCP/UDP Port Examples

**MSTP**





# Port Number

**MSTP**

- 0 – 255 are assigned to public applications
  - 80 is assigned for HTTP
- 256 – 1023 are assigned to “well known sockets”
  - 1752 is assigned for VTC
- 1024 and up are used to set up sessions
  - Randomly assigned



# User Datagram Protocol (UDP)

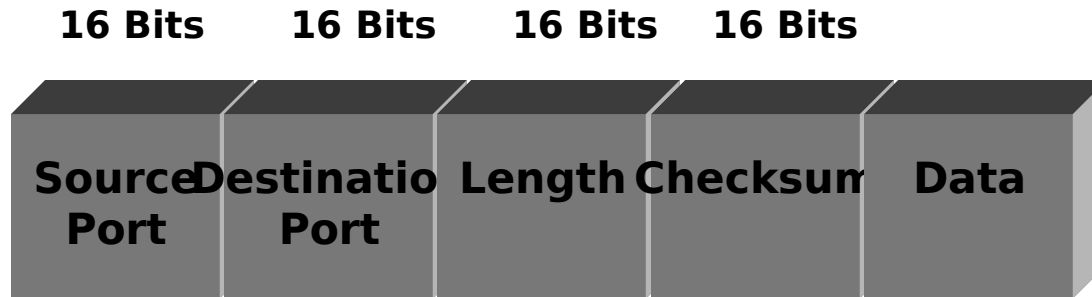
**MSTP**

- Used when all data fits in one packet
  - SNMP (Simple Network Management Protocol)
  - DNR (Domain Name Resolver)
  - NBT (NetBIOS over TCP/IP)
- Unreliable
  - No acknowledgment at this layer
- User data integrity
  - Adds header and computes checksum
- Why use UDP?
  - Lower overhead
  - Small amount of data for transmission
  - Less overhead to retransmit if data lost
  - Application entity has its own reliability built in



# UDP Header

**MSTP**



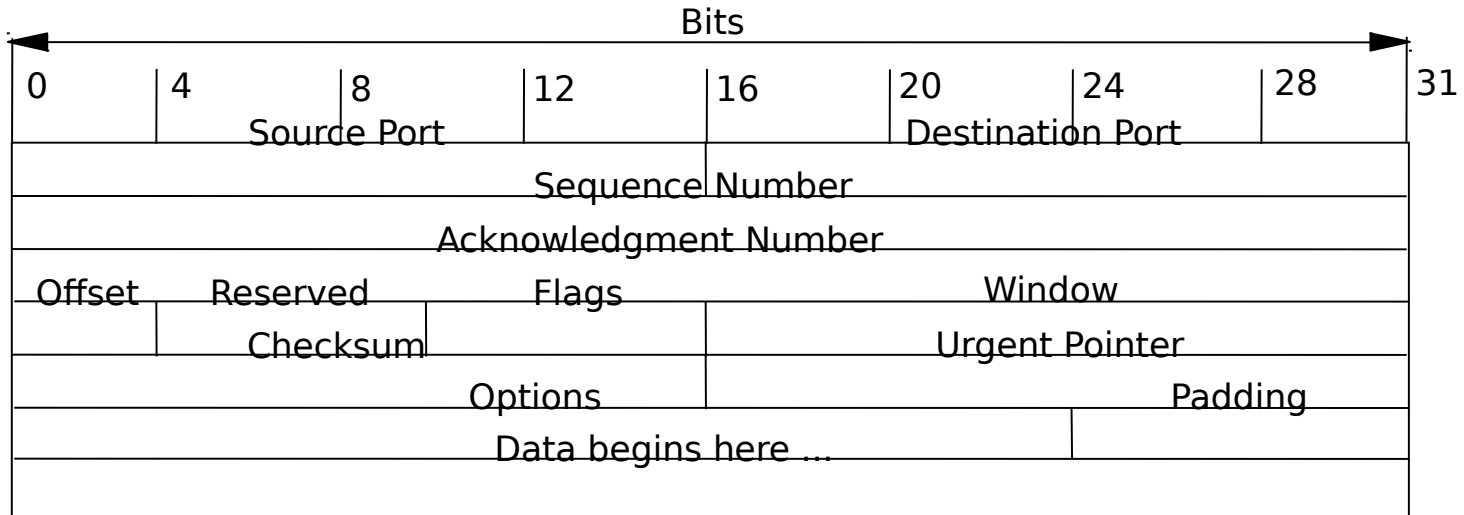




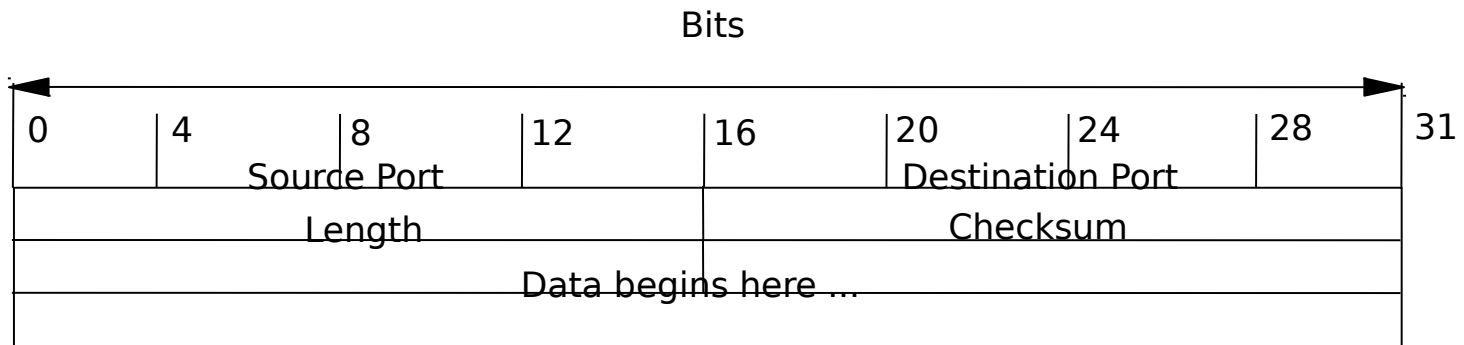
# TCP vs. UDP Header

**MSTP**

## TCP Header



## UDP Header





**Bring out the  
Sniffer**

No.	Status	Source Address	Dest Address	Summary	Len	Rel. Time	Delta Time	Abs. Time
22	#	[192.168.21.33]	[192.168.21.38]	Expert: ICMP Host Unreachable DLC: Ethertype=0800, size=70 bytes IP: D=[192.168.21.38] S=[192.168.21.33] LEN: ICMP: Destination unreachable (Host unreacha	70	0:00:12.033	0.000.735	
23		[192.168.21.38]	[192.168.20.10]	DLC: Ethertype=0800, size=60 bytes IP: D=[192.168.20.10] S=[192.168.21.38] LEN: TCP: D=23 S=1062 SYN SEQ=13074509 LEN=0 WIN=	60	0:00:12.341	0.308.605	
24		[192.168.20.10]	[192.168.21.38]	DLC: Ethertype=0800, size=60 bytes IP: D=[192.168.21.38] S=[192.168.20.10] LEN: TCP: D=1062 S=23 SYN ACK=13074510 SEQ=357988	60	0:00:12.359	0.017.653	
25		[192.168.21.38]	[192.168.20.10]	DLC: Ethertype=0800, size=60 bytes IP: D=[192.168.20.10] S=[192.168.21.38] LEN: TCP: D=23 S=1062 ACK=3579886763 WIN=8340	60	0:00:12.359	0.000.162	
26		[192.168.20.10]	[192.168.21.38]	DLC: Ethertype=0800, size=66 bytes IP: D=[192.168.21.38] S=[192.168.20.10] LEN: TCP: D=1062 S=23 ACK=13074510 SEQ=357988 Telnet: R PORT=1062 IAC Will Echo	66	0:00:12.378	0.018.964	
27		[192.168.21.38]	[192.168.20.10]	DLC: Ethertype=0800, size=60 bytes IP: D=[192.168.20.10] S=[192.168.21.38] LEN:	60	0:00:12.379	0.000.405	

IP: Version = 4, header length = 20 bytes
IP: Type of service = C0
IP: 110. .... = internetwork control
IP: ...0 .... = normal delay
IP: ....0... = normal throughput
IP: ....0... = normal reliability
IP: Total length = 52 bytes
IP: Identification = 1
IP: Flags = 0X
IP: .0... .... = may fragment
IP: ...0. .... = last fragment
IP: Fragment offset = 0 bytes
IP: Time to live = 254 seconds/hops
IP: Protocol = 6 (TCP)
IP: Header checksum = 1182 (correct)
IP: Source address = [192.168.20.10]
IP: Destination address = [192.168.21.38]

```

00000000: 00 00 86 35 6b 9f 00 e0 1e 3f b2 88 08 00 45 c0 ..!5k|..à.?²|..EA
00000010: 00 34 00 01 00 00 fe 06 11 82 c0 a8 14 0a c0 a8 .4....b...|À"...À"
00000020: 15 26 00 17 04 26 d5 60 bc a8 00 c7 80 4e 50 18 .&...&O'¼«ÇINP.
00000030: 10 20 ca a4 00 00 ff fb 01 ff fb 03 ff fd 18 ff .Ê...ÿà..ÿà..ÿÿ.ÿ
00000040: fd 1f

```



No.	Status	Source Address	Dest Address	Summary	Len	Rel. Time	Delta Time	Abs. Time
73		0030804FB30D	Bridge_Group_Ad	DLC: 802.3 size=38 bytes LLC: C D=42 S=42 UI BPDU: S: Pri=8000 Port=8019 Root: Pri=8000 Add:	64	0:00:28.002	0.157.129	
74	#	[192.168.21.38]	[138.156.24.250]	Expert: WINS No Response DLC: Ethertype=0800, size=110 bytes IP: D=[138.156.24.250] S=[192.168.21.38] IE UDP: D=137 S=137 LEN=76 WINS: C ID=35884 OP=REFRESH NAME=KORYNTAD<00	110	0:00:28.556	0.553.344	
75	#	[192.168.21.33]	[192.168.21.38]	Expert: ICMP Host Unreachable	70	0:00:28.557	0.000.816	

IP: Version = 4, header length = 20 bytes  
 IP: Type of service = 00  
 IP: 000. .... = routine  
 IP: ...0 .... = normal delay  
 IP: .... 0... = normal throughput  
 IP: .... .0.. = normal reliability  
 IP: Total length = 96 bytes  
 IP: Identification = 61458  
 IP: Flags = 0X  
 IP: .0... .... = may fragment  
 IP: ..0. .... = last fragment  
 IP: Fragment offset = 0 bytes  
 IP: Time to live = 128 seconds/hops  
 IP: Protocol = 17 (UDP)  
 IP: Header checksum = D115 (correct)  
 IP: Source address = [192.168.21.38]  
 IP: Destination address = [138.156.24.250]  
 IP: No options  
 IP:  
 UDP: ----- UDP Header -----  
 UDP:  
 UDP: Source port = 137 (NetBIOS-ns)

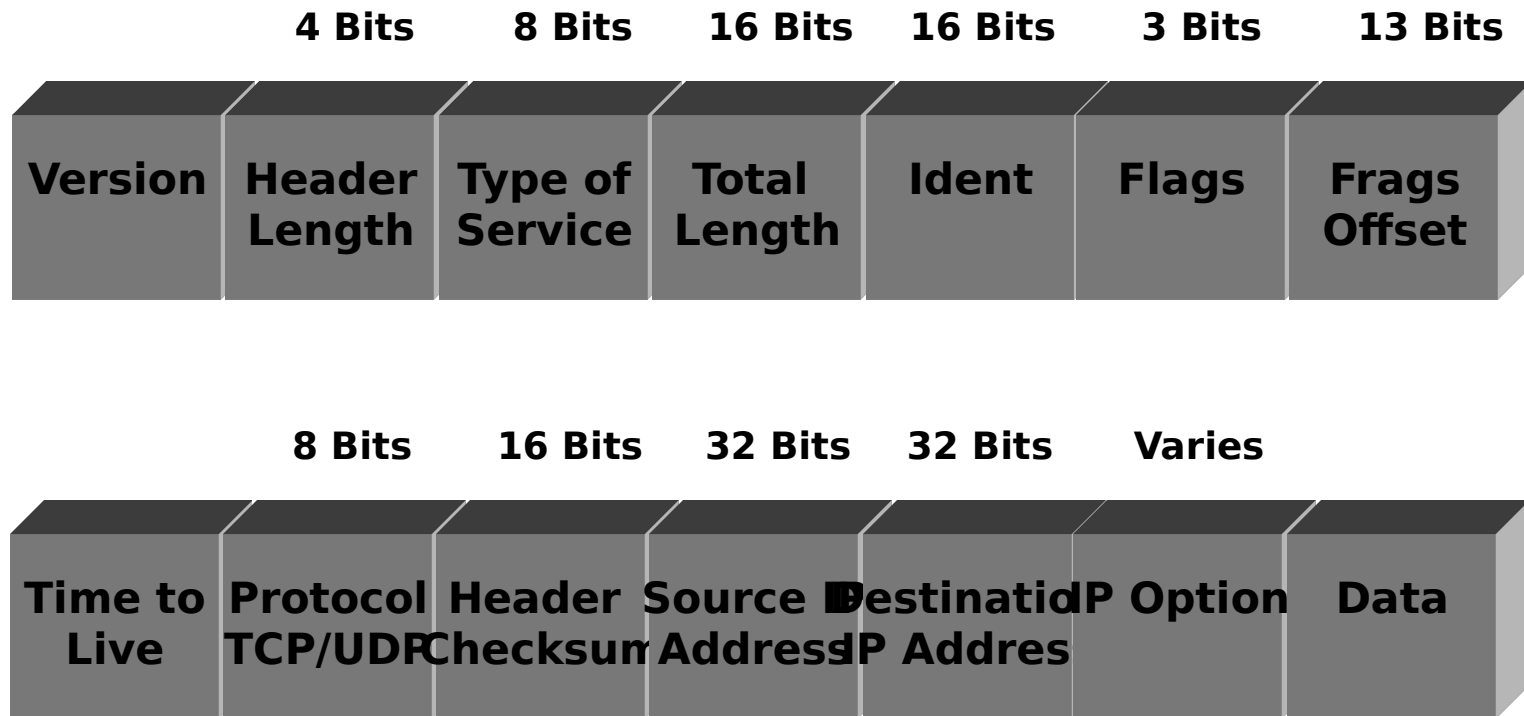
```

00000000: 00 e0 1e 3f b2 88 00 00 86 35 6b 9f 08 00 45 00 .à.?²...5k...E
00000010: 00 60 f0 12 00 00 80 11 d1 15 c0 a8 15 26 8a 9c .8....N.A"&||
00000020: 18 fa 00 89 00 89 00 4c be 62 8c 2c 40 00 00 01 .ú...L%b...@...
00000030: 00 00 00 00 00 01 20 45 4c 45 50 46 43 46 4a 45 .....ELEPFCEJE
00000040: 4f 46 45 45 42 45 45 43 41 43 41 43 41 43 41 43 OFEEBEECACACACAC
00000050: 41 43 41 43 41 41 41 00 00 20 00 01 c0 0c 00 20 ACACAAA...A...
00000060: 00 01 00 04 93 e0 00 06 60 00 c0 a8 15 26 .....à...A"&
  
```



# IP Header

**MSTP**



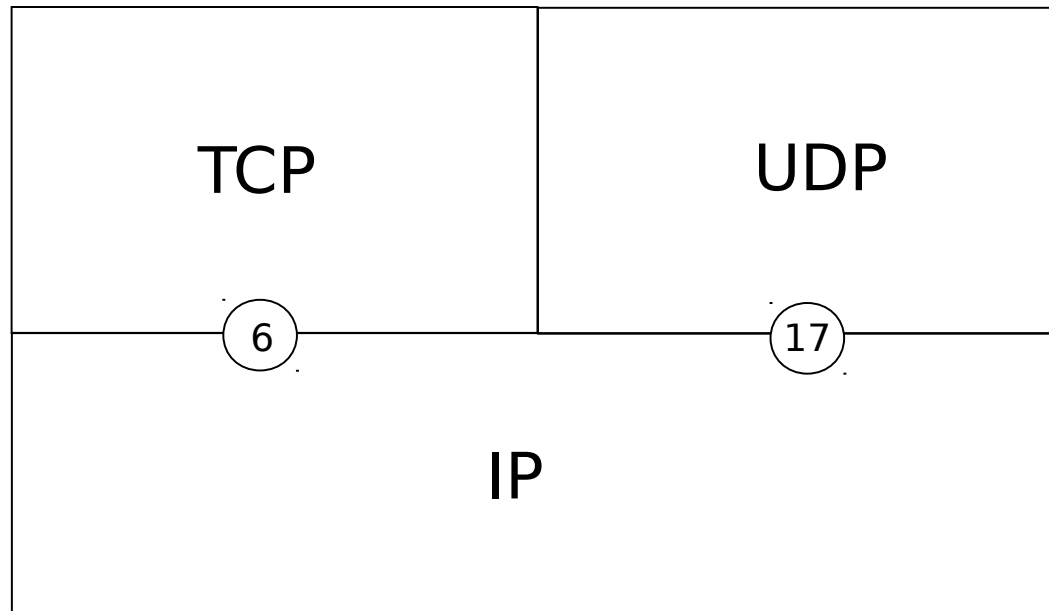


# IP Port has to be either

**MSTP**

**Transport  
Layer**

**Internet  
Layer**



Protocol  
Numbers



# Data Encapsulation

**MSTP**

Gen. Smith writes original letter on piece of paper .



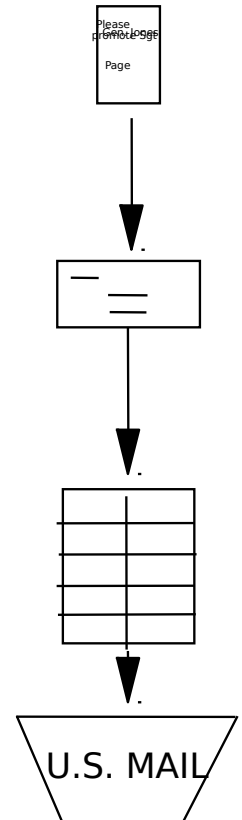
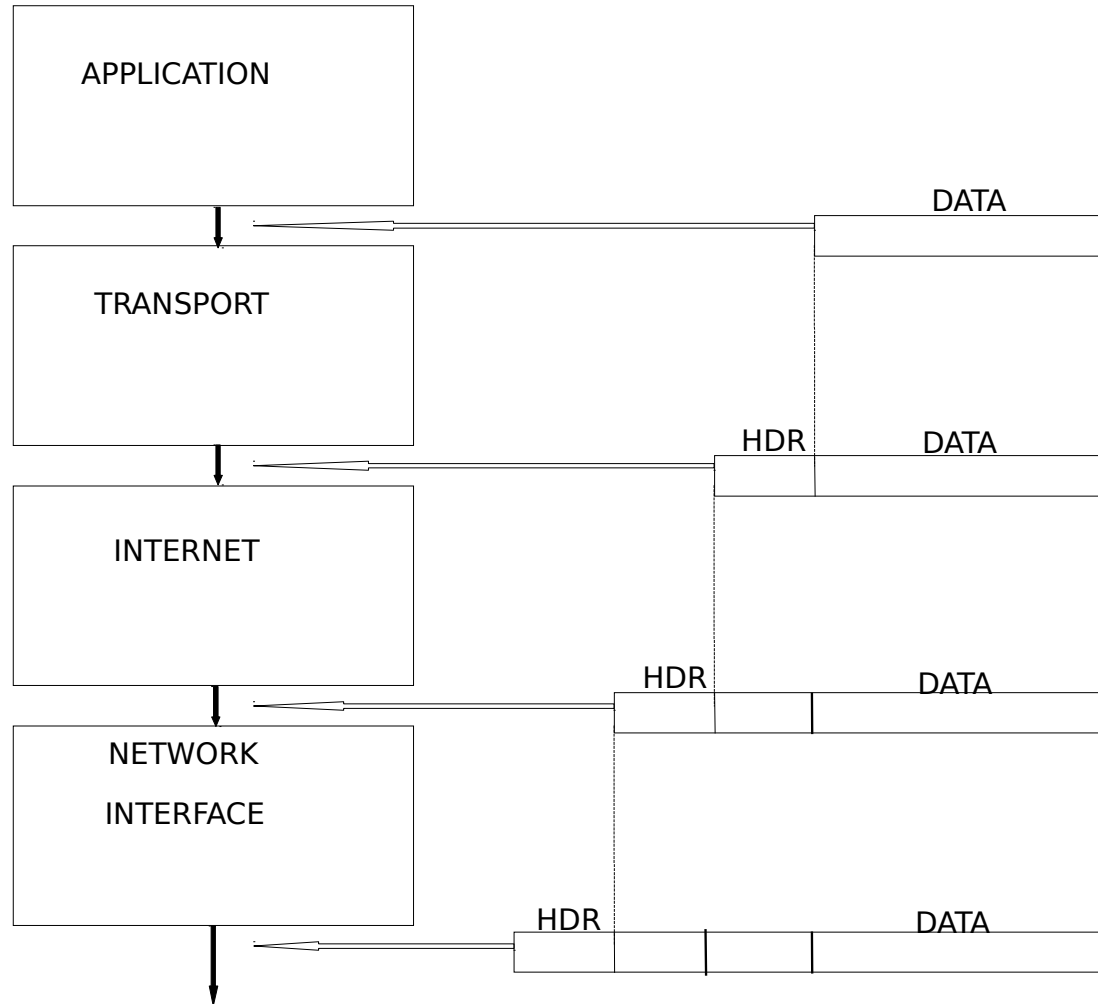
Maj. Bill puts the letter into an envelope.



Sgt York places the envelope into a guard mail package.



Pvt. Session throws the guard mail into his mail bag.





# How the computer sees the layers

## MSTP

APPLICATION LAYER

PROGRAM RUN BY USER

TRANSPORT LAYER

CONNECTION ORIENTED  
NON CONNECTION ORIENTED

INTERNET LAYER

NETWORK ADDRESS

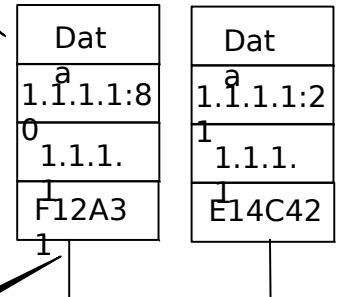
IP ADDRESS

NETWORK INTERFACE LAYER

HARDWARE ADDRESS

MAC ADDRESS

ETHERNET ADDRESS

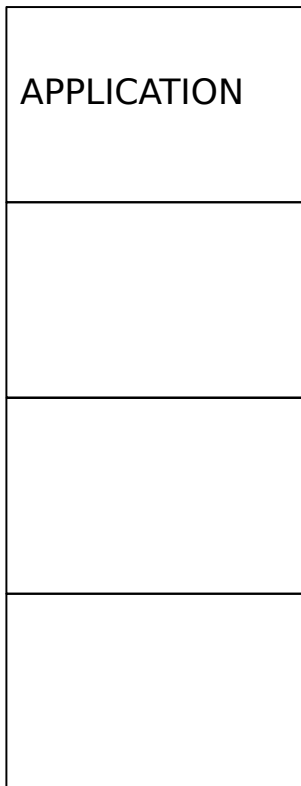






# Application Layer

**MSTP**



APPLICATION

HIGHEST LAYER

ONLY LAYER VISIBLE TO USERS

INTERACTS WITH TRANSPORT LAYER

DECIDES WHICH TRANSPORT PROTOCOL TO USE

COMMONLY KNOWN PROCESSES

TELNET

FTP

SMTP

DNS

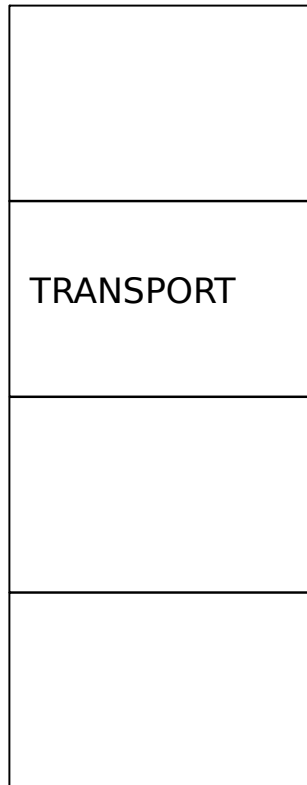
RIP

WWW



# Transport Layer

**MSTP**



RESPONSIBLE FOR HOST TO HOST DATA TRANSMISSION

KEEPS TRACK OF DATA RECEPTION & PASSES DATA TO APPLICATION ITSELF - TWO PROTOCOLS ARE :

TCP    RELIABLE

ACKNOWLEDGMENT

REGULATES DATA FLOW

ERROR DETECTION AND CORRECTION

UDP

UNRELIABLE

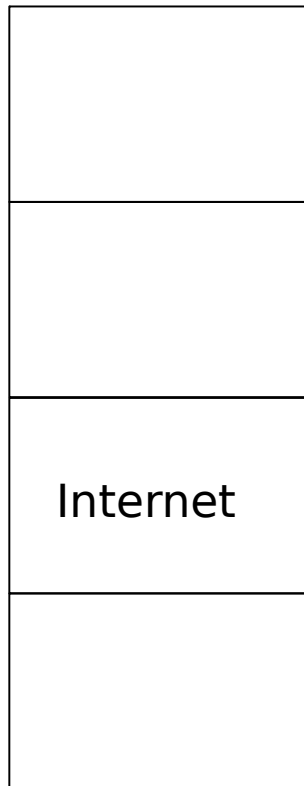
RELIES ON UPPER LAYERS FOR ERROR DETECTION

LOWER OVERHEAD



# Internet Layer

**MSTP**



Provides basic packet  
delivery  
Routing  
decisions  
Encapsulates packet in  
datagram  
Fragmentation and reassembly of  
packets  
Connectionless  
Unreliable



# Interface Layer

**MSTP**



- Accepts datagrams and transmits over media
- Node to node over single link
- Must know underlying transmission protocol
- One protocol for every network access protocol
- Data integrity
- Computes and checks checksums
- Encapsulation of datagrams
- Maps network addresses to hardware addresses
  - ARP
  - RARP



# Network Interface Layer Addressing

**MSTP**

- Hardware address
- MAC address
- Ethernet address
- 48 bits
- 05 23 33 20 00 f4
- 2 portions
  - Vendor code
    - First 6 characters
    - Assigned by IEEE
  - Unique hardware address
    - Last 6 characters
    - Assigned by vendor
- Actual address to which frames are sent



# Network Interface Card

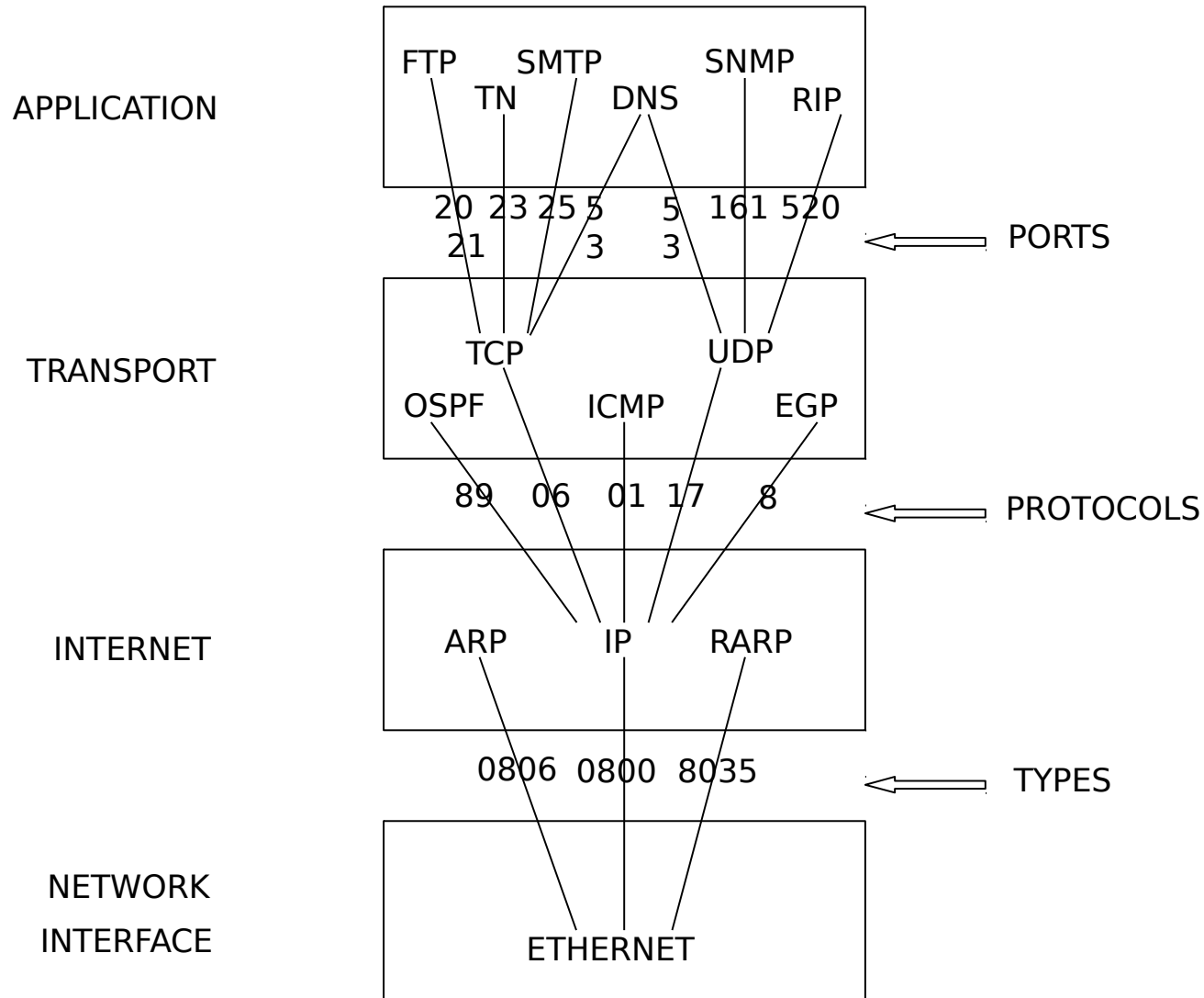
**MSTP**

- Network Interface Card (NIC) listens for:
  - It's hardware address
  - Broadcast address
  - Multicast address
- Decision Process
  - Ethernet NIC
    - Mine - Pass it to internet layer
    - Not Mine - Discard
  - Token Ring NIC
    - Mine - Pass it to internet layer
    - Not Mine - Regenerate and pass along



# Inter-layer Communication

**MSTP**





# Request for Comments (RFC)

**MSTP**

- The internet explains the open standard that makes up TCP/IP and related Internet protocols in Request for Comments (RFC's)
- RFC's are also written on many networking topics. Each new and received/replacement RFC is assigned a sequential number in the order that they are submitted.





# Address Resolution Protocol (ARP)

**MSTP**

- RFC 826
- Converts network address to hardware address
  - Deliver data from one host to another on same network
- Supported by majority of vendor's implementations
- Sender knows network address but not hardware address
  - Sender broadcasts ARP request to all hosts
  - All interfaces receive ARP request
    - If not mine, drop
    - If mine, reply to send
  - Sender caches hardware and network address
  - Sender sends data to recipient



# Reverse Address Resolution Protocol (RARP)

**MSTP**

- RFC 903 / 906
- Maps hardware addresses to network addresses
- Allows diskless clients to learn their own network addresses
- Workstation knows hardware address but not network address
  - Workstation broadcasts RARP request onto network
  - RARP server responds with network address
- Special chipset on NIC required
- RARP server must be available and



**MSTP**

# **IP Addressing & Classes**



# Table of Contents

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**MSTP**

- Binary Numbering
- Binary to Decimal
- IP Addressing
- Subnetting
- Questions
- Summarization



# Binary Numbering

**MSTP**

- Computer works using the binary numbering system.
- Recognizes two states, the **presence** of an electrical charge or the **absence** of an electrical charge. In other words, **on or off**.
- Binary numbering system is ideal for representing these two states
  - Consists of only two digits. (smallest – largest) 0 and 1
  - 0 represents the absence of an electrical charge or 'off'.
  - 1 represents the presence of an electrical charge or 'on'.
- **Bit = 1 digit (one or a zero)**
- **Byte = 8 bits**
- **Octet = Always 8 bits**



# Binary To Decimal

**MSTP**

**ONE** OCTET CAN BE BROKEN DOWN INTO

$$\begin{array}{r} 2^7 \quad 2^6 \quad 2^5 \quad 2^4 \quad 2^3 \quad 2^2 \quad 2^1 \quad 2^0 \\ \hline 128 \quad 64 \quad 32 \quad 16 \quad 8 \quad 4 \quad 2 \quad 1 \end{array}$$

$$2=2$$

$$2 \times 2 = 4$$

$$2 \times 2 \times 2 = 8$$

$$2 \times 2 \times 2 \times 2 = 16$$

$$2 \times 2 \times 2 \times 2 \times 2 = 32$$

$$2 \times 2 \times 2 \times 2 \times 2 \times 2 = 64$$

$$2 \times 2 \times 2 \times 2 \times 2 \times 2 \times 2 = 128$$

$$2^8$$



# MSTP

- | • | <u>128</u>     | <u>64</u> | <u>32</u> | <u>16</u> | <u>8</u> | <u>4</u> | <u>2</u> | <u>1</u> |
|---|----------------|-----------|-----------|-----------|----------|----------|----------|----------|
| • | 0              | 0         | 0         | 0         | 0        | 0        | 0        | 0        |
|   | 00000000 = 0   |           |           |           |          |          |          |          |
| • | 1              | 1         | 1         | 1         | 1        | 1        | 1        | 1        |
|   | 11111111 = 255 |           |           |           |          |          |          |          |
| • | 0              | 0         | 0         | 0         | 1        | 1        | 1        | 1        |
|   | 00001111 = 15  |           |           |           |          |          |          |          |
| • | 0              | 1         | 0         | 1         | 0        | 1        | 0        | 1        |
|   | 01010101 = 85  |           |           |           |          |          |          |          |



# Binary (Cont.)

**MSTP**

- 00000000 = 0
- 10000000 = 128
- 11000000 = 192
- 11100000 = 224
- 11110000 = 240
- 11111000 = 248
- 11111100 = 252
- 11111110 = 254
- 11111111 = 255



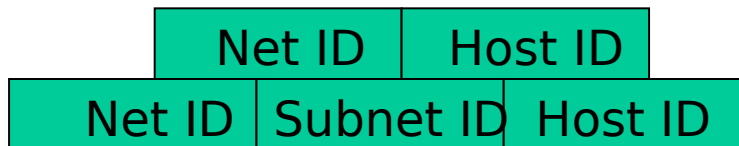




# IP Addressing

**MSTP**

- 32 BITS - 4 BYTES -
- MUST BE UNIQUE FOR EACH HOST IN NETWORK (8 BITS = 1 BYTE)
  - 192.168.20.10
  - 192.168.31.33
- 2 [3] PORTIONS
  - Network portion** - Relative to the class of IP. Identifies the network and is common to all devices attached to that network.
  - Host portion** - also relative to class as well as identifies a particular device attached to that network.
  - [Subnet portion]





# Addressing/Classes

**MSTP**

▲ XXX.XXX.XXX.XXX    **n = network    h=host address**  
▲ **192.156.2.169 (IPv4)**

● **Class A**    nnn.hhh.hhh.hhh    **1 - 126**  
-Only 126 networks, but 16,777,214 hosts apiece

▫ **127.0.0.1 = Local loop back address**

● **Class B**    nnn.nnn.hhh.hhh    **128 - 191**  
▫ 16,384 networks with 65,534 hosts apiece

● **Class C**    nnn.nnn.nnn.hhh    **192 - 223**  
▫ 2,097,152 networks with 254 hosts apiece

● **Class D - used for multicasting (audio/video)**

● **Class E - currently reserved / future**    **83**



# MSTP

84



# Decimal to Binary Example

**MSTP**

192.156.69.0 =

11000000.10011100.01000101.00000000

Class C

128	64	32	16	8	4	2	1	128	64	32	16	8	4	2	1	.	128	64	32	16	8	4	2	1
1	1	0	0	0	0	0	0	1	0	0	1	1	1	0	0	.	0	1	0	0	0	1	0	1
128+64 = 192								128+16+8+4 = 156								64+4+1 = 69								

Let Practice!!



# Reserved IP Addresses

## MSTP

- IP address for Hosts cannot have:

ALL 1's or ALL 0's (binary) in the NETWORK portion

OR

ALL 1's or ALL 0's (binary) in the HOST portion

- All 1's in the host portion of a target IP address signifies a Broadcast
- All 0's in the host portion of a IP address identify a subnet or network

- Network: 138.156.0.0 =  
10001010.10011110.00000000.00000000  
Broadcast: 138.156.255.255 =  
10001010.10011110.11111111.11111111  
Host: 138.156.100.100 =  
10001010.10011110.01100100.01100100



# IP Addresses

**MSTP**

- ASSIGNED by Node Site Coordinator
  - Address assignment planning
  - Node Site Coordinator
  - Draw out your network
  - Same "physical" net means same "IP Network"
  - Each "interface" has a "unique" IP address
  - "Don't" assign reserved addresses
- RECOMMENDATIONS
  - FIRST 10 addresses reserved for router interfaces
  - LAST address reserved for domain name



**MSTP**

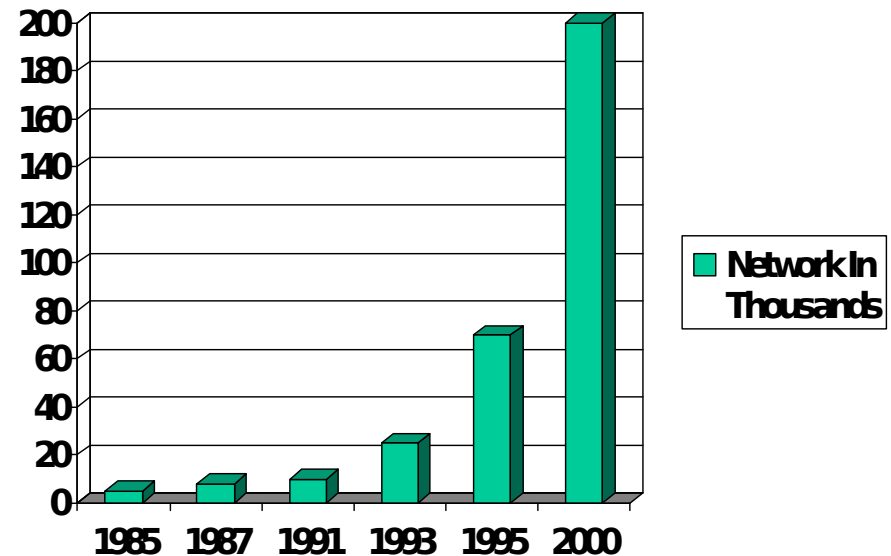
# **TCP/I P SUBNETTING**



# A little History to Why We Subnet

**MSTP**

- The growth of the Internet
- Depletion of IP Addresses - currently we have IPv4 which gives us 4,294,967,296 IP Address. ( $2^{32}$ )
- Administrator where not meeting current requirements based on the two-level classful IP address. Classful address were not be allocated efficiently.
- Growth caused Internet router's routing tables to increase







# Terminology

**MSTP**

- Address Mask – All network bits set to 1 and all host bits set to 0
- Subnet – A subnetwork of a major class A, B, C address space
- Subnet Mask – A mask longer than the standard address mask – determined by subnet scheme.



# IP Address Terminology

**MSTP**

- **NETWORK NUMBER**- When all host bits are turned off (0).
- **BROADCAST ADDRESS**- When all host bits are turned on (1).
- **HOST ADDRESS**- A unique IP address assigned to a workstation, interface or user, that is in between the network number and broadcast address.
- **SUBNET MASK**-Used to tell the machine what subnetting scheme is being implemented on the network. Found by turning all network bits on (1), including those host bits that have been given to the network side.
- **SUBNETTING** - Dividing up an entire Class network by sacrificing original host (H) bits to the network (N).



# SUBNETTING

**MSTP**

- WHAT IS IT?  
Divides host (H) portion into smaller networks
- WHY?  
Stops wasting network numbers
- WHO?  
Node site coordinator
- WHAT DETERMINES?  
Number of different physical networks and  
number of hosts



# Subnetting

**MSTP**

- When you borrow bits from the main network address's host section, TCP/IP must be told which bits of the host section are borrowed to be used as the network address.

- We use a subnet mask to define the number of bits used to create additional networks.

- **Remember** - the more bits used to define the mask, the fewer the



# Default Subnet Mask

**MSTP**

Your network has a subnet mask even if it doesn't have

**CLASS A DEFAULT = 255.0.0.0**

**11111111.00000000.00000000.00000000**

**CLASS B DEFAULT = 255.255.0.0**

**11111111.11111111.00000000.00000000**

**CLASS C DEFAULT = 255.255.255.0**

**11111111.11111111.11111111.00000000**



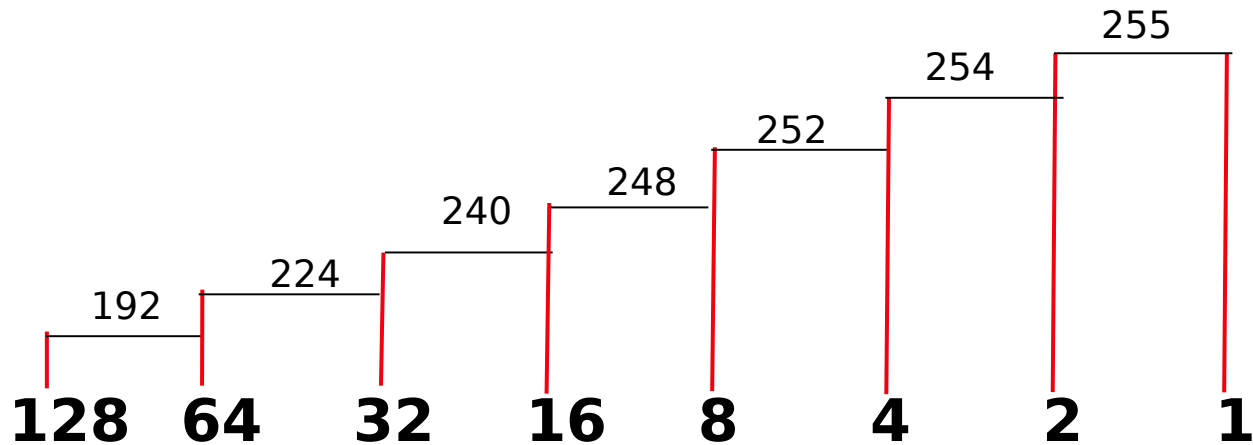
# MSTP

95



# Subnet Bit Chart

**MSTP**



\* SUBNET BITS COME FROM THE **HIGHEST-ORDER BITS** TO THE LOW ORDER BITS OF THE **HOST FIELD**



# How Subnetting Works

**MSTP**

N | H  
128 64 32 16 8 4 2  
1 MASK =  
128+64+32+16  
240

16 - (2) # hosts  
per network  
- all 0's network  
- all 1's broadcast  
leaves 14 per net

N				H					
128	64	32	16	8	4	2	1		
0	0	0	1	0	0	0	0	16	Network
0	0	0	1	0	0	0	1	17	1st host
0	0	0		1	0	0	1	0	18
0	0	0		1	0	0	1	1	19
:	:	:		:	:	:	:	:	
0	0	0		1	1	1	1	1	31
Broadcast									
0	0	1		0	0	0	0	0	32





# Determining Subnet Mask

**MSTP**

192.156.69.0 = 11000000.10011100.01000101.00000000 Class C  
 N.N.N.H

143.211.0.0 = 10001111.11010011.00000000.00000000 Class B

Subnet Mask is the address with every network bit turned on. This tells the router that you want to use some Host bits as network (subnet) bits.

192.156.69.0 = 11000000.10011100.01000101.00000000 Class C  
 N.N.N.H

Subnet Mask = 11111111.11111111.11111111.00000000 255.255.255.0

with 4 bit = 11111111.11111111.11111111.11110000  
 255.255.255.240

143.211.0.0 = 10001111.11010011.00000000.00000000 Class B  
 N.N.H.H

Subnet Mask = 11111111.11111111.00000000.00000000 255.255.0.0

with 8 bit = 11111111.11111111.11111111.00000000



# Subnetting Reference Charts

**MSTP**

## CLASS B

# BITS	SUBNET MASK	# SUBNETS	# HOSTS
2	255.255.192.0	2	16382
3	255.255.224.0	6	8190
4	255.255.240.0	14	4094
5	255.255.248.0	30	2046
6	255.255.252.0	62	1022
7	255.255.254.0	126	510
8	255.255.255.0	254	254
9	255.255.255.128	510	126
10	255.255.255.192	1022	62
11	255.255.255.224	2046	30
12	255.255.255.240	4094	14
13	255.255.255.248	8190	6
14	255.255.255.252	16382	2



# Subnetting Reference Charts

**MSTP**

## CLASS C

# BITS	SUBNET MASK	# SUBNETS	# HOSTS
2	255.255.255.192	4	62
3	255.255.255.224	6	30
4	255.255.255.240	14	14
5	255.255.255.248	30	6
6	255.255.255.252	62	2



# Steps in Subnetting

**MSTP**

1. Write Out the Subnet Mask
2. Answer What You Know
3. Write Out in Binary
4. Apply Logical And (or Anding)
5. Turn on all the host bits

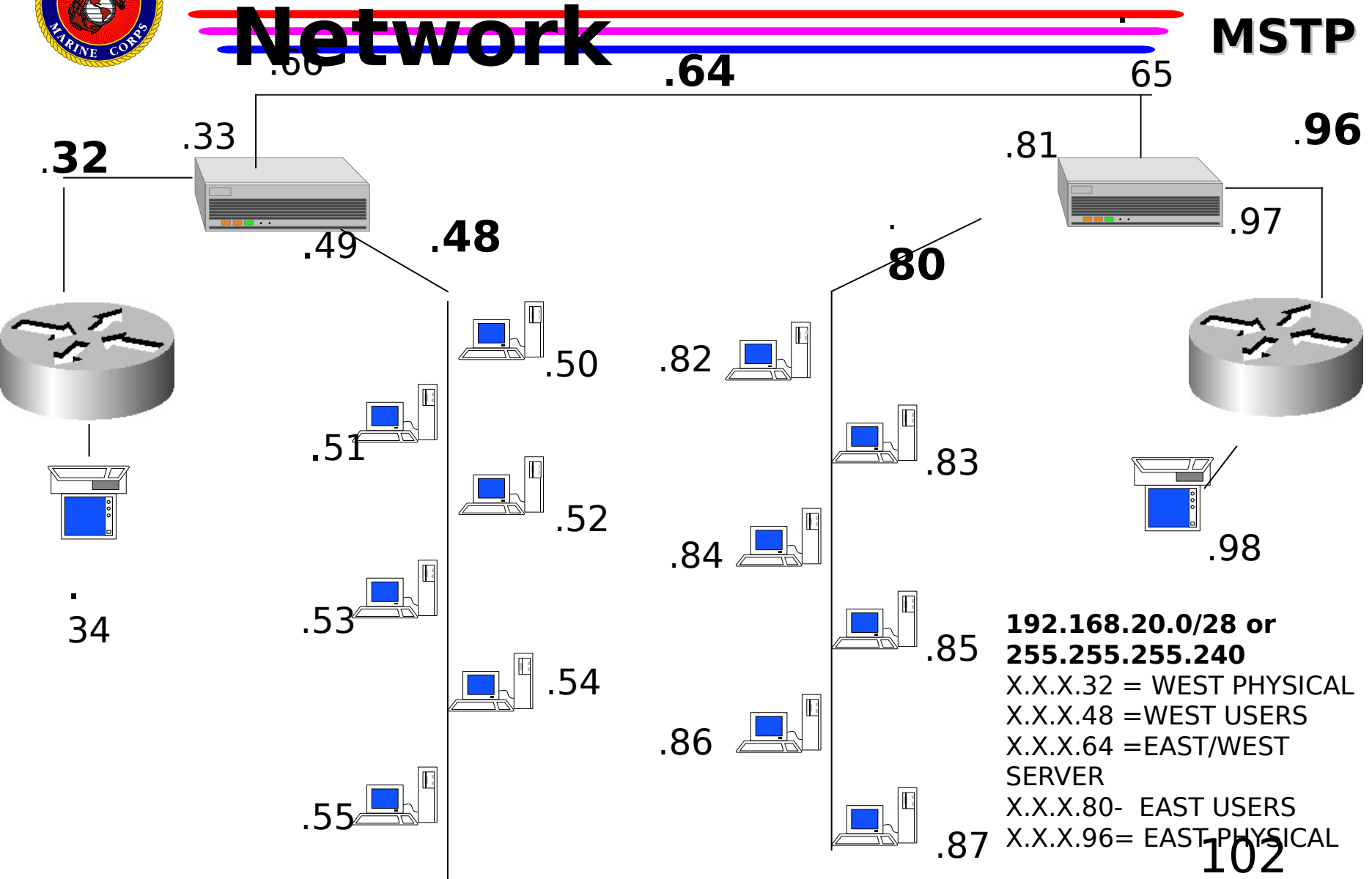
**192.168.25.45**  
**/ 27**  
 192.168.25.45  
 255.255.255.224  
 192.168.25.??  
 128 64 32 16 8 4  
 2 1  
 1 1 1 0 0 0  
 0 0  
 0 0 1 0 1 1  
 0 1  
 0 0 1 0 0  
 0 0 0  
 0 0 1 1 0 1  
 1 1 1



# Sample Subnetted

## Network

**MSTP**





# Going Beyond The Octet

**MSTP**

N										H									
32768	16384	8192	4096	2048	1024	512	256			128	64	32	16	8	4	2	1		
128	64	32	16	8	4	2	1	●		128	64	32	16	8	4	2	1		
0	0	0	0	0	0	0	0			0	1	0	0	0	0	0	0	0.64 Network	
0	0	0	0	0	0	0	0			0	1	0	0	0	0	0	1	0.65 1st host	
0	0	0	0	0	0	0	0			0	1	0	0	0	0	1	0	0.66 2nd Host	
0	0	0	0	0	0	0	0			0	1	0	0	0	0	0	0		
0	0	0	0	0	0	0	0			0	1	1	1	1	1	1	1	0.127 Broadcast	
0	0	0	0	0	0	0	0			1	0	0	0	0	0	0	0	0.128 Network	
0	0	0	0	0	0	0	0			1	0	0	0	0	0	0	1	0.129 1st host	
32768	16384	8192	4096	2048	1024	512	256	●		128	64	32	16	8	4	2	1		
128	64	32	16	8	4	2	1			0	0	0	0	0	0	0	0	4.0 Network	
0	0	0	0	0	0	1	0			0	0	0	0	0	0	0	1	4.1 1st host	
0	0	0	0	0	0	1	1			1	1	1	1	1	1	1	1	7.255 Broadcast	
0	0	0	0	1	0	0	0			0	0	0	0	0	0	0	0	8.0 Network	

**10 Bit**  
1022 Nets  
62 Hosts

**6 Bit**  
62 Nets  
1022 Hosts



# Finding A Host's Network

**MSTP**

**N**

**H**

Host: 192.156.69.78 = 11000000.10011100.01000101.01001110

Subnet Mask 4 bit = 11111111.11111111.11111111.11111111

Host Network ID # = 11000000.10011100.01000101.01001110  
which is not all 1's or all 0's.

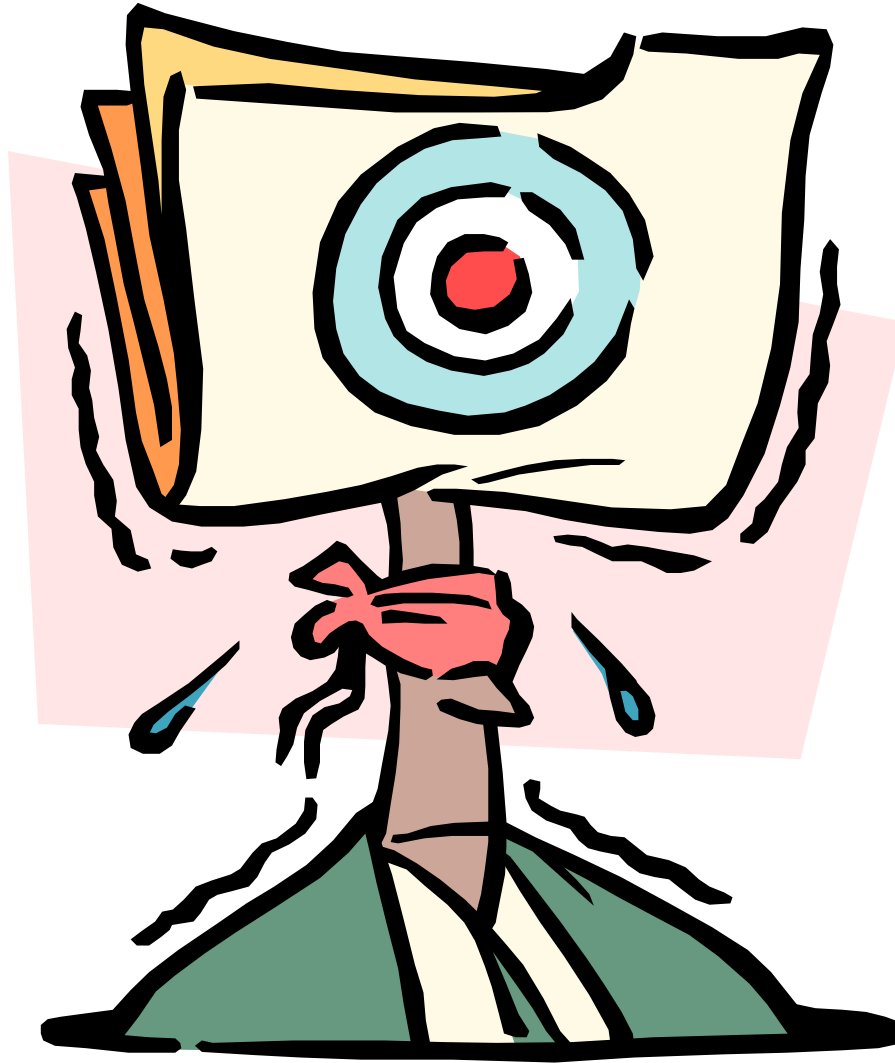
Host bits = 1110 which is not all 1's or all 0's so it is legal, it is the

14th host on the .64 network.



# Is there another way?

**MSTP**







# Five Questions

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---

**MSTP**

- 1. How many subnets?
- 2. How many hosts per subnet?
- 3. What are the subnets?
- 4. What are the valid hosts in each subnet?
- 5. What is the broadcast address of each subnet?



## Begin to answer by...

**MSTP**

1. Determine how many networks you need.
2. Find out how many hosts are required for each network (use the highest number of hosts).
3. Choose the subnetting scheme that will best support all networks (leave room for growth).
4. Assign network numbers.
5. Assign unique addresses to



# Five Answers

**MSTP**

- 1.  $2^n$  = Amount of subnets.
- 2.  $2^n - 2$  = Amount of hosts per subnet.
- 3.  $256 - \text{Subnet mask} = \text{Base number}$ .
- 4. Valid hosts are the numbers between the subnets, minus all 0's and all 1's.
- 5. Broadcast address is all 1's or the number before the next subnet.



# Prefix Routing

**MSTP**

- Means by which the Internet identifies the portion of the 32-bit TCP/IP address
  - /27      255.255.255.224
  - /26      255.255.255.192
  - /25      255.255.255.128
  - /24      255.255.255.0
  - /23      255.255.254.0



# Discontiguous Addressing

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## MSTP

- Two networks of the same classful networks are separated by a different network address.
- When using RIP or IGRP, you must use the default-router command.



# Summarization

**MSTP**

- Allows contiguous networks to be grouped together and advertised as one large network
- Also known as supernetting



# Any Questions

**MSTP**

